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Francesca Poggiolesi
**A Cut-Free Simple Sequent Calculus
for Modal Logic S5**

Abstract

In this paper we present a simple sequent calculus for the modal propositional logic S5. We prove that this sequent calculus is theoremwise equivalent to the Hilbert-style system S5, that it is contraction-free and cut-free, and finally that it is decidable. All results are proved in a purely syntactic way.

1 Introduction

Amongst the many normal systems of modal propositional logic, one of the most important and well-known is doubtlessly S5.

When considered from the point of view of Kripke semantics, S5 is quite a peculiar system since it can be described in two different but equivalent ways¹. The first one specifies the properties that the accessibility relation between worlds of a Kripke frame should satisfy: S5 is indeed sound and complete with respect to the class of reflexive, transitive and symmetric frames (or, equivalently, in the class of reflexive and euclidean frames). A second and easier way to study S5 semantically exploits Kripke frames where the accessibility relation is absent²: S5 is indeed sound and complete with respect to the class of frames which are just non-empty sets of worlds.

This second way is evidently simpler and it would be useful and interesting to reflect this simplicity at the syntactic level, within, for example, a Gentzen system. Unfortunately this has not yet been achieved: indeed, it turns out to be quite a challenge to give a sequent calculus for S5. The efforts in this direction are numerous and each of them presents some difficulties. The most common fault of the sequent calculi proposed so far is not being cut-free (Blumey and Humberstone (1991), Matsumoto and Ohnishi (1959)), or, in any case, not satisfying the subformula property (Negri (2005), Sato (1980), and Wansing (1994)). Other Gentzen calculi are syntactically impure because they use explicit semantic parameters (Braüner (2000), Cerrato (1993), Mints (1997)). Finally there exist Gentzen systems for S5 which are quite laborious since they treat S5 as a system whose accessibility relation satisfies several conditions (Brünnler (2006), Došen (1985), Indrzejczak (1997)).

The best solutions are probably those offered by Avron (1996) and Restall (2006): both these solutions (i) use hypersequents, (ii) try to reflect the simpler

¹For a detailed explanation of this point see Restall (2006).

²From now on, we will call this kind of frame an S5 Kripke frame.

way S5 can be described semantically. Since the main goal of this paper is to present a new sequent calculus for S5 which shares with Avron and Restall's solutions the points (i) and (ii), it is worth explaining in detail how our approach differs from these two. The major differences are the following. Firstly, in our calculus all the structural rules are (height-preserving) admissible (even the contraction-rules), and the logical and modal rules are height-preserving invertible. Secondly, the rules which compose our calculus are different from the rules of the other calculi. Thirdly, we prove the decidability theorem in a purely syntactic way. Finally, as has been shown in Poggiolesi (2007), our calculus is closely related to the so called tree-hypersequents method (Poggiolesi (2008)) by means of which we can obtain Gentzen systems for several modal logics. Like Restall, but in contrast to Avron, our rules are explicit³. Finally we point out that Restall is the only one to propose a proofnets interpretation of his sequent calculus.

2 The calculus CSS5_s

We define the modal propositional language \mathcal{L}^\square in the following way:

- atoms: p_0, p_1, \dots
- logical constant: \square ,
- connectives: \neg, \vee .

The other connectives, the modal operator \diamond , and the formulas of the modal language \mathcal{L}^\square are defined as usual.

Syntactic Conventions:

- α, β, \dots : formulas,
- M, N, P, Q, \dots : finite multisets of formulas,
- Γ, Δ, \dots : classical sequents.
- G, H, \dots : hypersequents.

As we will deal with hypersequents, we remind the reader what a hypersequent is.

Definition 2.1. A hypersequent is a syntactic object of the form:

$$M_1 \Rightarrow N_1 | M_2 \Rightarrow N_2 | \dots | M_n \Rightarrow N_n$$

where $M_i \Rightarrow N_i$ ($i = 1, \dots, n$) are classical sequents.

³Intuitively, the rules of a sequent calculus are explicit when they can be presented in such a way that the symbol they introduce does not appear in their premise(s). For a definition of the explicitness property see Wansing (1994).

Definition 2.2. The intended interpretation of a hypersequent is defined inductively in the following way:

- $(M \Rightarrow N)^\tau := \bigwedge M \rightarrow \bigvee N$,
- $(\Gamma_1 | \Gamma_2 | \dots | \Gamma_n)^\tau := \Box \Gamma_1^\tau \vee \Box \Gamma_2^\tau \vee \dots \vee \Box \Gamma_n^\tau$

Given the definition and the interpretation of the notion of hypersequent, it should be clear that a hypersequent is just a *multiset* of classical sequents, which is to say that the order of the sequents in a hypersequents is not important.

The postulates of the calculus $CSS5_s$ ⁴ are:

Initials Hypersequents

$$G \mid M, p \Rightarrow N, p$$

Propositional Rules

$$\frac{G \mid M \Rightarrow N, \alpha}{G \mid \neg \alpha, M \Rightarrow N} \neg A \qquad \frac{G \mid \alpha, M \Rightarrow N}{G \mid M \Rightarrow N, \neg \alpha} \neg K$$

$$\frac{G \mid \alpha, \beta, M \Rightarrow N}{G \mid \alpha \wedge \beta, M \Rightarrow N} \wedge A \qquad \frac{G \mid M \Rightarrow N, \alpha \quad G \mid M \Rightarrow N, \beta}{G \mid M \Rightarrow N, \alpha \wedge \beta} \wedge K$$

Modal Rules

$$\frac{G \mid \alpha, \Box \alpha, M \Rightarrow N}{G \mid \Box \alpha, M \Rightarrow N} \Box A_1 \qquad \frac{G \mid M \Rightarrow N \mid \Rightarrow \alpha}{G \mid M \Rightarrow N, \Box \alpha} \Box K$$

$$\frac{G \mid \Box \alpha, M \Rightarrow N \mid \alpha, P \Rightarrow Q}{G \mid \Box \alpha, M \Rightarrow N \mid P \Rightarrow Q} \Box A_2$$

Let us make two remarks on the modal rules. The first one only concerns the rules $(\Box A_i)$ ($i = 1, 2$). The repetition of the principal formula $\Box \alpha$ in the premise of each of these rules only serves to make the rules invertible. This is analogous to the repetition of the formula $(\forall x)(\alpha(x))$ in the premise of the rule which

⁴The name $CSS5_s$ stands for: *simple* version of the calculus $CSS5$. Indeed in Poggiolesi (2006) we have introduced another sequent calculus for $S5$, called precisely $CSS5$, which reflects the more complicated way $S5$ can be described semantically that was mentioned in the introduction.

introduces the symbol \forall on the left side of the sequent in some versions of the sequent calculus for first-order classical logic.

The second remark concerns the three modal rules. It is easy to informally understand these rules if we compare the hypersequent to a S5 Kripke frame, and the sequents which compose the hypersequent to different worlds of the S5 Kripke frame. In this perspective, the rule $(\Box K)$ says, if read bottom up, that, if the formula $\Box\alpha$ is false at a world x , then we can create a new world y where the formula α is false; on the other hand, the rules $(\Box A_i)$ tell us, if read bottom up and considered together, that, if the formula $\Box\alpha$ is true at a world x , then the formula α is true in *any* world of the frame.

3 Admissibility of the Structural Rules

In this section we will show which structural rules are admissible in the calculus $CSS5_s$. Moreover, in order to show that the two rules of contraction are height-preserving admissible, we will show that all the logical and modal rules are height-preserving invertible. The proof of the admissibility of the cut-rule will be shown in the fifth section.

Definition 3.1. We associate to each proof d in $CSS5_s$ a natural number $h(d)$ (the height). Intuitively, the height corresponds to the length of the longest branch in a tree-proof d , minus one. We define $h(d)$ by induction on the construction of d .

$$\begin{aligned}
d &\equiv \quad G \mid M \Rightarrow N : & h(d) &= 0 \\
&\quad \vdots \\
d &\equiv \frac{G' \mid M' \Rightarrow N'}{G \mid M \Rightarrow N} \mathcal{R} : & h(d) &= h(d_1) + 1 \\
&\quad \vdots \\
d &\equiv \frac{\begin{array}{c} \vdots \\ G \mid M' \Rightarrow N' \end{array} \quad \begin{array}{c} \vdots \\ G \mid M'' \Rightarrow N'' \end{array}}{G \mid M \Rightarrow N} \mathcal{R} : & h(d) &= \max (h(d_1) + 1, h(d_2) + 1)
\end{aligned}$$

Definition 3.2. $d \vdash^n G$ means that d is a proof of G in $CSS5_s$, with $h(d) \leq n$. We write $\vdash^{(n)} G$, or just $\langle^n \rangle G$, for: “there exists a proof d such that $d \vdash^n G$.”

Definition 3.3. Let G be a hypersequent and G' be the result of the application of a certain rule \mathcal{R} on G . We say that this rule \mathcal{R} is *height-preserving admissible* when:

$$d \vdash^n G \Rightarrow \exists d' (d' \vdash^n G')$$

We call a rule, \mathcal{R} , which transforms a hypersequent G into a hypersequent G' , *admissible* when:

$$d \vdash^n G \Rightarrow \exists d' (d' \vdash G')$$

Lemma 3.4. *Hypersequents of the form $G|\alpha, M \Rightarrow N, \alpha$, with α an arbitrary modal formula, are derivable in $CSS5_s$.*

Proof. By straightforward induction on α . \square

Lemma 3.5. *The rule of merge:*

$$\frac{G | M \Rightarrow N | P \Rightarrow Q}{G | M, P \Rightarrow N, Q} \text{merge}$$

is height-preserving admissible in $CSS5_s$.

Proof. By induction on the height of the derivation of the premise.

If the premise is an initial hypersequent, then so is the conclusion.

If the premise is inferred by a logical rule, then the inference is clearly preserved. We will give an example using the logical rule ($\neg K$):

$$\frac{\langle n-1 \rangle G | \alpha, M \Rightarrow N | P \Rightarrow Q}{\langle n \rangle G | M \Rightarrow N, \neg \alpha | P \Rightarrow Q} \neg K \quad \rightsquigarrow^5$$

$$\frac{\langle n-1 \rangle G | \alpha, M, P \Rightarrow N, Q}{\langle n \rangle G | M, P \Rightarrow N, Q, \neg \alpha} \neg K$$

If the premise is inferred by the modal rule ($\Box K$), this is clearly preserved:

$$\frac{\langle n-1 \rangle G | M \Rightarrow N | P \Rightarrow Q | \Rightarrow \alpha}{\langle n \rangle G | M \Rightarrow N, \Box \alpha | P \Rightarrow Q} \Box K \quad \rightsquigarrow$$

$$\frac{\langle n-1 \rangle G | M, P \Rightarrow N, Q | \Rightarrow \alpha}{\langle n \rangle G | M, P \Rightarrow N, Q, \Box \alpha} \Box K$$

⁵The symbol \rightsquigarrow means: the premise of the right side is obtained by induction hypothesis on the premise of the left side.

If the premise is inferred by the modal rule ($\Box A_1$), this is clearly preserved:

$$\frac{\langle^{n-1}\rangle G \mid \alpha, \Box\alpha, M \Rightarrow N \mid P \Rightarrow Q}{\langle^n\rangle G \mid \Box\alpha, M \Rightarrow N \mid P \Rightarrow Q} \Box A_1 \quad \rightsquigarrow$$

$$\frac{\langle^{n-1}\rangle G \mid \Box\alpha, \alpha, M, P \Rightarrow N, Q}{\langle^n\rangle G \mid \Box\alpha, M, P \Rightarrow N, Q} \Box A_1$$

If the premise is inferred by the modal rule ($\Box A_2$), there are two significant cases to analyze: the one where the rule ($\Box A_2$) has been applied between the two sequents $M \Rightarrow N$ and $P \Rightarrow Q$; the one where the rule ($\Box A_2$) has been applied between a third sequent, let us call it Γ , and $M \Rightarrow N$ (or, equivalently, $P \Rightarrow Q$). The two situations are analogous, therefore we only analyze the first one in detail:

$$\frac{\langle^{n-1}\rangle G \mid \Box\alpha, M \Rightarrow N \mid \alpha, P \Rightarrow Q}{\langle^n\rangle G \mid \Box\alpha, M \Rightarrow N \mid P \Rightarrow Q} \Box A_2 \quad \rightsquigarrow$$

$$\frac{\langle^{n-1}\rangle G \mid \Box\alpha, \alpha, M, P \Rightarrow N, Q}{\langle^n\rangle G \mid \Box\alpha, M, P \Rightarrow N, Q} \Box A_1$$

⊠

Lemma 3.6. *The rule of external weakening:*

$$\frac{G}{G \mid M \Rightarrow N}^{EW}$$

is height-preserving admissible in $CSS5_s$.

Proof. By straightforward induction on the height of the derivation of the premise. ⊠

Lemma 3.7. *The rule of internal weakening:*

$$\frac{G \mid M \Rightarrow N}{G \mid M, P \Rightarrow N, Q}^{IW}$$

is height-preserving admissible in $CSS5_s$.

Proof. It follows by the height-preserving admissibility of the two rules of merge and external weakening. \square

Lemma 3.8. *All the logical and modal rules of $CSS5_s$ are height-preserving invertible.*

Proof. The proof proceeds by induction on the height of the derivation of the premise of the rule considered. The cases of logical rules are dealt with in the classical way. The only differences - the fact that we are dealing with hypersequents and the cases where the last applied rule is one of the rules $(\Box A_i)$ or $(\Box K)$ - are dealt with easily.

The rules $(\Box A_i)$ are trivially height-preserving invertible since both their premises are obtained by weakening from their respective conclusions, and weakening is height-preserving admissible.

We show in detail the invertibility of the rule $(\Box K)$. If $G \mid M \Rightarrow N, \Box \alpha$ is an initial hypersequent, then so is $G \mid M \Rightarrow N \mid \Rightarrow \alpha$. If $G \mid M \Rightarrow N, \Box \alpha$ is obtained by a logical rule \mathcal{R} , we apply the inductive hypothesis on the premise(s), $G' \mid M' \Rightarrow N', \Box \alpha$ ($G'' \mid M'' \Rightarrow N'', \Box \alpha$) and we obtain derivation(s), of height $n - 1$, of $G' \mid M' \Rightarrow N' \mid \Rightarrow \alpha$ ($G'' \mid M'' \Rightarrow N'' \mid \Rightarrow \alpha$). By applying the rule \mathcal{R} , we obtain a derivation of height n of $G \mid M \Rightarrow N \mid \Rightarrow \alpha$. If $G \mid M \Rightarrow N, \Box \alpha$ is of the form $G \mid \Box \beta, M' \Rightarrow N, \Box \alpha$, then it may have been obtained by the two modal rules $(\Box A_i)$. Since the procedure is the same in both cases, we can just analyze the one of the rule $(\Box A_1)$ and the other can be dealt with analogously. We apply the inductive hypothesis on $G \mid \Box \beta, \beta, M' \Rightarrow N, \Box \alpha$ and we obtain a derivation of height $n - 1$ of $G \mid \Box \beta, \beta, M' \Rightarrow N \mid \Rightarrow \alpha$. By applying the rule $(\Box A_1)$, we obtain a derivation of height n of $G \mid \Box \beta, M' \Rightarrow N \mid \Rightarrow \alpha$. If $G \mid M \Rightarrow N, \Box \alpha$ is obtained by the modal rule $(\Box K)$ in which $\Box \alpha$ is not the principal formula, then this case can be treated analogously to the one of $(\Box A_i)$. Finally, if $G \mid M \Rightarrow N, \Box \alpha$ is obtained by the modal rule $(\Box K)$ and $\Box \alpha$ is the principal formula, the premise of the last step gives the conclusion. \square

Lemma 3.9. *The rules of contraction:*

$$\frac{G \mid \alpha, \alpha, M \Rightarrow N}{G \mid \alpha, M \Rightarrow N} \text{CA} \qquad \frac{G \mid M \Rightarrow N, \alpha, \alpha}{G \mid M \Rightarrow N, \alpha} \text{CK}$$

are height-preserving admissible in $CSS5_s$.

Proof. By induction on the height of the derivation of the premise $G \mid M \Rightarrow N, \alpha, \alpha$ ($G \mid \alpha, \alpha, M \Rightarrow N$). We analyze only the case of the rule CK . The case of the rule CA is symmetric.

If $G \mid M \Rightarrow N, \alpha, \alpha$ is an initial hypersequent, so is $G \mid M \Rightarrow N, \alpha$.

If $G \mid M \Rightarrow N, \alpha, \alpha$ is the conclusion of a rule \mathcal{R} which does not have either of the two occurrences of the formula α as principal, we apply the inductive

hypothesis on the premise(s) $G' \mid M' \Rightarrow N', \alpha, \alpha$ ($G'' \mid M'' \Rightarrow N'', \alpha, \alpha$), obtaining derivation(s) of height $n - 1$ of $G' \mid M' \Rightarrow N', \alpha$ ($G'' \mid M'' \Rightarrow N'', \alpha$). By applying the rule \mathcal{R} we obtain a derivation of height n of $G \mid M \Rightarrow N, \alpha$.

$G \mid M \Rightarrow N, \alpha, \alpha$ is the conclusion of a logical or modal rule and one of the two occurrences of the formula α is principal. Hence the last rule used in the proof of $G \mid M \Rightarrow N, \alpha, \alpha$ is a K -rule and we have to analyze the following three cases: $\neg K$, $\wedge K$, $\Box K$.

$[\neg K]$:

$$\frac{\langle^{n-1}\rangle G \mid \beta, M \Rightarrow N, \neg\beta}{\langle^n\rangle G \mid M \Rightarrow N, \neg\beta, \neg\beta} \neg K \quad \dashrightarrow^6 \quad \frac{\langle^{n-1}\rangle G \mid \beta, \beta, M \Rightarrow N}{\langle^{n-1}\rangle G \mid \beta, M \Rightarrow N} i.h.}{\langle^n\rangle G \mid M \Rightarrow N, \neg\beta} \neg K$$

$[\wedge K]$:

$$\frac{\langle^{n-1}\rangle G \mid M \Rightarrow N, \beta, \beta \wedge \gamma \quad \langle^{n-1}\rangle G \mid M \Rightarrow N, \gamma, \beta \wedge \gamma}{\langle^n\rangle G \mid M \Rightarrow N, \beta \wedge \gamma, \beta \wedge \gamma} \wedge K \quad \dashrightarrow$$

$$\frac{\frac{\langle^{n-1}\rangle G \mid M \Rightarrow N, \beta, \beta}{\langle^{n-1}\rangle G \mid M \Rightarrow N, \beta} i.h. \quad \frac{\langle^{n-1}\rangle G \mid M \Rightarrow N, \gamma, \gamma}{\langle^{n-1}\rangle G \mid M \Rightarrow N, \gamma} i.h.}{\langle^n\rangle G \mid M \Rightarrow N, \beta \wedge \gamma} \wedge K$$

$[\Box K]$:

$$\frac{\langle^{n-1}\rangle G \mid M \Rightarrow N, \Box\beta \mid \Rightarrow \beta}{\langle^n\rangle G \mid M \Rightarrow N, \Box\beta, \Box\beta} \Box K \quad \dashrightarrow$$

$$\frac{\frac{\langle^{n-1}\rangle G \mid M \Rightarrow N \mid \Rightarrow \beta \mid \Rightarrow \beta}{\langle^{n-1}\rangle G \mid M \Rightarrow N \mid \Rightarrow \beta, \beta} merge}{\langle^{n-1}\rangle G \mid M \Rightarrow N \mid \Rightarrow \beta} i.h.}{\langle^n\rangle G \mid M \Rightarrow N, \Box\beta} \Box K$$

⊠

⁶The symbol \dashrightarrow means: the premise of the right side is obtained by an application of Lemma 3.8 on the premise of the left side.

4 The Adequateness Theorem

In this section we briefly prove that the sequent calculus $CSS5_s$ proves exactly the same formulas as its corresponding Hilbert-style system $S5$.

Theorem 4.1.

(i) If $\vdash \alpha$ in $S5$, then $\vdash \Rightarrow \alpha$ in $CSS5_s$.

(ii) If $\vdash G$ in $CSS5_s$, then $\vdash (G)^\tau$ in $S5$.

Proof. By induction on the height of proofs in $S5$ and $CSS5_s$, respectively. In order to further acquaint the reader with the calculus $CSS5_s$ we verify (i). The classical axioms and the modus ponens are proved as usual, we just present the proof of the axiom T , the axiom 4, the axiom B and the axiom 5.

$CSS5_s \vdash \Rightarrow \Box\alpha \rightarrow \alpha$

$$\frac{\frac{\Box\alpha, \alpha \Rightarrow \alpha}{\Box\alpha \Rightarrow \alpha}}{\Rightarrow \Box\alpha \rightarrow \alpha}$$

$CSS5_s \vdash \Rightarrow \Box\alpha \rightarrow \Box\Box\alpha$

$$\frac{\frac{\frac{\frac{\Box\alpha \Rightarrow | \Rightarrow |\alpha \Rightarrow \alpha}{\Box\alpha \Rightarrow | \Rightarrow | \Rightarrow \alpha}}{\Box\alpha \Rightarrow | \Rightarrow \Box\alpha}}{\Box\alpha \Rightarrow \Box\Box\alpha}}{\Rightarrow \Box\alpha \rightarrow \Box\Box\alpha}}$$

$CSS5_s \vdash \Rightarrow \alpha \rightarrow \Box\neg\Box\neg\alpha$

$$\frac{\frac{\frac{\frac{\alpha \Rightarrow \alpha | \Box\neg\alpha \Rightarrow}{\neg\alpha, \alpha \Rightarrow | \Box\neg\alpha \Rightarrow}}{\alpha \Rightarrow | \Box\neg\alpha \Rightarrow}}{\alpha \Rightarrow | \Rightarrow \neg\Box\neg\alpha}}{\frac{\alpha \Rightarrow \Box\neg\Box\neg\alpha}{\Rightarrow \alpha \rightarrow \Box\neg\Box\neg\alpha}}$$

$CSS5_s \vdash \Rightarrow \neg\Box\neg\alpha \rightarrow \Box\neg\Box\neg\alpha$

$$\begin{array}{c}
\Rightarrow |\Box\neg\alpha \Rightarrow | \alpha \Rightarrow \alpha \\
\hline
\Rightarrow |\Box\neg\alpha \Rightarrow | \Rightarrow \neg\alpha, \alpha \\
\hline
\Rightarrow |\Box\neg\alpha \Rightarrow | \neg\alpha \Rightarrow \neg\alpha \\
\hline
\Rightarrow |\Box\neg\alpha \Rightarrow | \Rightarrow \neg\alpha \\
\hline
\Rightarrow | \Rightarrow \neg\Box\neg\alpha | \Rightarrow \neg\alpha \\
\hline
\Rightarrow \Box\neg\alpha | \Rightarrow \neg\Box\neg\alpha \\
\hline
\Rightarrow \Box\neg\alpha, \Box\neg\Box\neg\alpha \\
\hline
\neg\Box\neg\alpha \Rightarrow \Box\neg\Box\neg\alpha \\
\hline
\Rightarrow \neg\Box\neg\alpha \Rightarrow \Box\neg\Box\neg\alpha
\end{array}$$

⊠

5 Cut-elimination Theorem

In this section we prove the cut-elimination theorem for the calculus $CSS5_g$.

Theorem 5.1. *Let $G \mid M \Rightarrow N, \alpha$ and $G' \mid \alpha, P \Rightarrow Q$ be two hypersequents, we want to prove that:*

$$\text{if } \frac{\begin{array}{c} \vdots \\ d_1 \{ G \mid M \Rightarrow N, \alpha \quad G' \mid \alpha, P \Rightarrow Q \}^{d_2} \\ \vdots \end{array}}{G \mid G' \mid M, P \Rightarrow N, Q} \text{cut}_\alpha$$

and d_1 and d_2 do not contain any other application of the cut-rule, then we can construct a proof of $G \mid G' \mid M, P \Rightarrow N, Q$ without any application of the cut-rule.

Proof. This is proved by induction on the complexity of the cut-formula, which is the number (≥ 0) of the occurrences of logical symbols in cut formula α , with subinduction on the sum of the heights of the derivations of the premises of the cut. We will distinguish cases by the last rule applied on the left premise.

Case 1. $G \mid M \Rightarrow N, \alpha$ is an initial hypersequent. Then either the conclusion is also an initial hypersequent or the cut can be replaced by various applications of the internal and external weakening rules on $G' \mid \alpha, P \Rightarrow Q$.

Case 2. $G \mid M \Rightarrow N, \alpha$ is inferred by a rule \mathcal{R} in which α is not principal. The reduction is carried out in the standard way by induction on the sum of the heights of the derivations of the premises of cut.

Case 3. $G \mid M \Rightarrow N, \alpha$ is inferred by a rule \mathcal{R} in which α is principal. We distinguish two subcases: in one subcase \mathcal{R} is a propositional rule, in the other \mathcal{R} is a modal rule.

Case 3.1. As an example, we consider the case where the rule before $G \mid M \Rightarrow N$, α is $\neg K$, we have:

$$\frac{\frac{G \mid \beta, M \Rightarrow N}{G \mid M \Rightarrow N, \neg\beta} \neg K \quad \frac{G' \mid \neg\beta, P \Rightarrow Q}{G \mid G' \mid M, P \Rightarrow N, Q} \text{cut}_{\neg\beta}}{\vdots}$$

By applying lemma 3.8 on $G' \mid \neg\beta, P \Rightarrow Q$, we obtain $G' \mid P \Rightarrow Q, \beta$. Therefore we can replace the previous cut with the following one which is eliminable by induction on the complexity of the cut formula:

$$\frac{G' \mid P \Rightarrow Q, \beta \quad G \mid \beta, M \Rightarrow N}{G \mid G' \mid M, P \Rightarrow N, Q} \text{cut}_{\beta}$$

Case 3.2. \mathcal{R} is $\Box K$ and $\alpha \equiv \Box\beta$. We have the following situation:

$$\frac{\frac{G \mid M \Rightarrow N \mid \Rightarrow \beta}{G \mid M \Rightarrow N, \Box\beta} \Box K \quad \frac{G' \mid \Box\beta, P \Rightarrow Q}{G \mid G' \mid M, P \Rightarrow N, Q} \text{cut}_{\Box\beta}}{\vdots}$$

We have to consider the last rule \mathcal{R}' of d_2 . If there is no rule \mathcal{R}' which introduces $G' \mid \Box\beta, P \Rightarrow Q$ because $G' \mid \Box\beta, P \Rightarrow Q$ is an initial hypersequent, then we can solve the case as in 1. If \mathcal{R}' is a rule in which $\Box\beta$ is not principal, we solve the case as in 2. The only problematic cases are thus those in which \mathcal{R}' is one of the rules $(\Box A_i)$. Since the procedure is the same in both cases, we need only analyze the case of the rule $(\Box A_2)$; the other case is dealt with analogously.

$$\frac{\frac{G \mid M \Rightarrow N \mid \Rightarrow \beta}{G \mid M \Rightarrow N, \Box\beta} \Box K \quad \frac{G' \mid \Box\beta, P \Rightarrow Q \mid \beta, Z \Rightarrow W}{G' \mid \Box\beta, P \Rightarrow Q \mid Z \Rightarrow W} \Box A_2}{G \mid G' \mid M, P \Rightarrow N, Q \mid Z \Rightarrow W} \text{cut}_{\Box\beta}$$

We reduce to:

$$\frac{\frac{G \mid M \Rightarrow N, \Box\beta \quad G' \mid \Box\beta, P \Rightarrow Q \mid \beta, Z \Rightarrow W}{G \mid G' \mid M, P \Rightarrow N, Q \mid \beta, Z \Rightarrow W} \text{cut}_{\Box\beta}}{\frac{G \mid M \Rightarrow N \mid \Rightarrow \beta \quad G \mid G' \mid M, P \Rightarrow N, Q \mid \beta, Z \Rightarrow W}{G \mid G' \mid M, P \Rightarrow N, Q \mid Z \Rightarrow W} \text{cut}_{\beta}} \text{merge}^* C^*$$

where the first cut is eliminable by induction on the sum of the heights of the derivations of the premises of cut and the second cut is eliminable by induction on the complexity of cut formula. \boxtimes

6 Decidability

In this section we prove that the calculus $CSS5_s$ is decidable, which is to say there is an algorithm that, given any hypersequent G , determines whether G is provable in $CSS5_s$ or not.

First of all let us observe that our calculus satisfies the subformula property since: (i) the cut rule is admissible in it (see theorem 5.1), (ii) in each of its rules all the formulas which occur in the premise(s) are subformulas of the formulas which occur in the conclusion. Moreover it can also be shown that the contraction rules are admissible (see lemma 3.9). It would therefore seem that any source of potentially non-terminating proof search had been cut off. Unfortunately it is not so because of the repetition of the principal formula in each of the rules $(\Box A_i)$. In order to avoid this problem and prove that our calculus is decidable, we shall obtain a bound on the number of applications of the rules $(\Box A_i)$.

For this goal, let us start by only taking into account *minimal derivations*, which is to say, derivations where shortenings are not possible. Then we prove, by means of the following lemmas and their corollaries, that in minimal derivations it is enough to apply the rule $(\Box A_1)$ only once on any given pair of principal formulas and the rule $(\Box A_2)$ only once on any given pair of sequents. This technique is inspired by the one used in Negri (2005).

Lemma 6.1. *The rule $(\Box A_1)$ permutes down with respect to the rules $(\neg A)$, $(\neg K)$, $(\wedge A)$, $(\wedge K)$, $(\Box A_2)$ and $(\Box K)$.*

Proof. Let us first of all consider the permutation with one-premise logical rules, which is straightforward. Taking, as an example, the case of the rule $(\neg K)$, we have:

$$\frac{\frac{G \mid \beta, \alpha, \Box \alpha, M \Rightarrow N}{G \mid \beta, \Box \alpha, M \Rightarrow N} \Box A_1}{G \mid \Box \alpha, M \Rightarrow N, \neg \beta} \neg K$$

$$\downarrow$$

$$\frac{\frac{G \mid \beta, \alpha, \Box \alpha, M \Rightarrow N}{G \mid \alpha, \Box \alpha, M \Rightarrow N, \neg \beta} \neg K}{G \mid \Box \alpha, M \Rightarrow N, \neg \beta} \Box A_1$$

Let us now consider the permutation with the two premises-rule ($\wedge K$). We have the following derivation:

$$\begin{array}{c}
\frac{G \mid \alpha, \Box\alpha, M \Rightarrow N, \beta}{G \mid \Box\alpha, M \Rightarrow N, \beta} \Box A_1 \quad \vdots \quad G \mid \Box\alpha, M \Rightarrow N, \gamma}{G \mid \Box\alpha, M \Rightarrow N, \beta \wedge \gamma} \wedge K \\
\downarrow \\
\frac{\vdots \quad G \mid \Box\alpha, M \Rightarrow N, \gamma}{G \mid \alpha, \Box\alpha, M \Rightarrow N, \gamma} IW \quad \frac{G \mid \alpha, \Box\alpha, M \Rightarrow N, \beta \wedge \gamma}{G \mid \Box\alpha, M \Rightarrow N, \beta \wedge \gamma} \Box A_1}{G \mid \alpha, \Box\alpha, M \Rightarrow N, \beta \wedge \gamma} \wedge K
\end{array}$$

Let us remark that the transformation of the first derivation into the second one is done by means of an application of the height-preserving admissible rule of internal weakening (IW).

Finally we show the permutation in case of the rule ($\Box K$):

$$\begin{array}{c}
\frac{G \mid \alpha, \Box\alpha, M \Rightarrow N \mid \Rightarrow \beta}{G \mid \Box\alpha, M \Rightarrow N \mid \Rightarrow \beta} \Box A_1}{G \mid \Box\alpha, M \Rightarrow N, \Box\beta} \Box K \\
\downarrow \\
\frac{G \mid \alpha, \Box\alpha, M \Rightarrow N \mid \Rightarrow \beta}{G \mid \alpha, \Box\alpha, M \Rightarrow N, \Box\beta} \Box K}{G \mid \Box\alpha, M \Rightarrow N, \Box\beta} \Box A_1
\end{array}$$

□

Lemma 6.2. *The rule ($\Box A_2$) permutes down with respect to the rules ($\neg A$), ($\neg K$), ($\wedge A$), ($\wedge K$), ($\Box A_1$). It also permutes with the instances of ($\Box K$) when the principal formula α of its premise is not active in the sequent where the principal formula of the premise of $\Box K$ occurs.*

Proof. The cases where one permutes with the one-premise propositional rules or with the two-premises rule ($\wedge K$) can be dealt with analogously to the corresponding ones of the previous lemma. We show the permutation in case of the rule ($\Box K$). We underline that this case is constrained by the hypothesis that the principal formula α of the premise of the rule ($\Box A_2$) is not active in the same

sequent where the principal formula of the premise of the rule ($\Box K$) occurs. By taking in account this restriction, the permutation is straightforward:

$$\frac{\frac{G \mid \Box\alpha, M \Rightarrow N \mid \alpha, P \Rightarrow Q \mid \Rightarrow \beta}{G \mid \Box\alpha, M \Rightarrow N \mid P \Rightarrow Q \mid \Rightarrow \beta} \Box A_2}{G \mid \Box\alpha, M \Rightarrow N, \Box\beta \mid P \Rightarrow Q} \Box K$$

↓

$$\frac{\frac{G \mid \Box\alpha, M \Rightarrow N \mid \alpha, P \Rightarrow Q \mid \Rightarrow \beta}{G \mid \Box\alpha, M \Rightarrow N, \Box\beta \mid \alpha, P \Rightarrow Q} \Box K}{G \mid \Box\alpha, M \Rightarrow N, \Box\beta \mid P \Rightarrow Q} \Box A_2$$

⊠

Corollary 6.3. *In a minimal derivation in $CSS5_s$, the rule ($\Box A_1$) cannot be applied more than once on the same pair of principal formulas of any branch.*

Proof. Let us suppose to have a minimal derivation where the rule ($\Box A_1$) has been applied twice on the same pair of sequents:

$$\frac{G' \mid \alpha, \Box\alpha, M' \Rightarrow N'}{G' \mid \Box\alpha, M' \Rightarrow N'} \Box A_1$$

⋮

$$\frac{G \mid \alpha, \Box\alpha, M \Rightarrow N}{G \mid \Box\alpha, M \Rightarrow N} \Box A_1$$

By permuting down ($\Box A_1$) with respect to the the steps in the dotted part of the derivation, we obtain a derivation of the same height ending with:

$$\frac{\frac{G \mid \alpha, \alpha, \Box\alpha, M \Rightarrow N}{G \mid \alpha, \Box\alpha, M \Rightarrow N} \Box A_1}{G \mid \Box\alpha, M \Rightarrow N} \Box A_1$$

By applying the height-preserving admissible rule (CA) on the two occurrences of the formula α in place of the upper ($\Box A_1$), we obtain a shorter derivation, contrary to the assumption of minimality. ⊠

Corollary 6.4. *In a minimal derivation in $CSS5_s$, the rule $(\Box A_2)$ cannot be applied more than once on the same pair of sequents of any branch.*

Proof. Consider a minimal derivation where the rule $(\Box A_2)$ has been applied twice on the same pair of sequents:

$$\frac{G' \mid \Box\alpha, M' \Rightarrow N' \mid \alpha, P' \Rightarrow Q'}{G' \mid \Box\alpha, M' \Rightarrow N' \mid P' \Rightarrow Q'} \Box A_2$$

.

.

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$$\frac{G \mid \Box\alpha, M \Rightarrow N \mid \alpha, P \Rightarrow Q}{G \mid \Box\alpha, M \Rightarrow N \mid P \Rightarrow Q} \Box A_2$$

By permuting down $(\Box A_2)$ with respect to the steps in the dotted part of the derivation, we obtain a derivation of the same height ending with:

$$\frac{\frac{G \mid \Box\alpha, M \Rightarrow N \mid \alpha, \alpha, P \Rightarrow Q}{G \mid \Box\alpha, M \Rightarrow N \mid \alpha, P \Rightarrow Q} \Box A_2}{G \mid \Box\alpha, M \Rightarrow N \mid P \Rightarrow Q} \Box A_2$$

By applying the height-preserving admissible rule (CA) on the two occurrences of the formula α in place of the upper $(\Box A_2)$, we obtain a shorter derivation, contrary to the assumption of minimality.

Finally we underline that if the principal formulas of the premise of the upper application of the rule $(\Box A_2)$ were active in the sequent where the principal formula of the premise of $\Box K$ occurs, then *that sequent would disappear* and therefore we would not find it in the premise of the lower application of the rule $(\Box A_2)$. The conclusion is that the restriction of lemma 6.2 is not limitative for the proof of this corollary. \square

Now we prove that the modal logic $S5$ is decidable by showing effective bounds on proof search in the calculus $CSS5_s$.

Theorem 6.5. *The calculus $CSS5_s$ allows terminating proof search.*

Proof. Place an hypersequent G , for which we are looking for a proof search, at the root of the procedure. Apply first the propositional rules and then the modal rules. The propositional rules reduce the complexity of the hypersequent. The rule $(\Box K)$ removes the modal constant \Box and adds a new sequent, each of the rules $(\Box A_i)$ increases the complexity. However, by the corollary 6.3, the rule $(\Box A_1)$ cannot be applied more than once on the same pair of principal formulas,

while, by the corollary 6.4, the rule $(\Box A_2)$ cannot be applied more than once on the same pair of sequents. Therefore the number of applications of the two rules $(\Box A_1)$ and $(\Box A_2)$ is bounded, respectively, by the number of \Box 's occurring in the negative part (see definition below) of the hypersequent to prove, and by the number of the sequents which may appear in the derivation. This derivation, in its turn, is bounded by the number of sequents belonging to the hypersequent to prove, and the sequents which can be introduced by applications of the rule $(\Box K)$.

We finally explain how to calculate explicit bounds. First of all define the *negative* and *positive* parts of the hypersequent $M_1 \Rightarrow N_1 | \dots | M_n \Rightarrow N_n$, as the negative and positive parts of each of the following conjuncts and disjuncts:

$$\bigwedge M_1 \rightarrow \bigvee N_1, \dots, \bigwedge M_n \rightarrow \bigvee N_n$$

For any given hypersequent G , let $n(\Box)$ be the number of \Box 's in the negative part of the hypersequent G , and $p(\Box)$ be the number of \Box 's in the positive part of the hypersequent G .

The number of applications of the rule $(\Box A_1)$ in a minimal derivation is bounded by:

$$n(\Box)$$

In the case where the root-hypersequent is just a sequent, the number of applications of the rule $(\Box A_2)$ in a minimal derivation is bounded by:

$$n(\Box) \cdot p(\Box)$$

In the case where the root-hypersequent is a hypersequent and s is the number of sequents which occurs in it, the number of applications of the rule $(\Box A_2)$ in a minimal derivation is bounded by:

$$n(\Box) \cdot (p(\Box) + s)$$

□

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Francesca Poggiolesi,
 Department of Philosophy,
 University of Florence,
 IHPST,
 University of Paris 1,

Via Bolognese 52,
50139, Florence, Italy.
email: francesca.poggiolesi@unifi.it

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