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Emergence on decreasing sandpile models*

Submitted version: complete proofs of Lemmas 1, 2, 3, 4, 5,
Theorems 4, 5, and definition of Muller automata in appendix.

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Abstract. Sand is a proper instance for the study of natural algorithmic phenomena. Idealized square/cubic sand grains moving according to “simple” local toppling rules may exhibit surprisingly “complex” global behaviors. In this paper we explore the language made by words corresponding to fixed points reached by iterating a toppling rule starting from a finite stack of sand grains in one dimension. Using arguments from linear algebra, we give a constructive proof that for all decreasing sandpile rules the language of fixed points is accepted by a finite (Muller) automaton. The analysis is completed with a combinatorial study of cases where the emergence of precise regular patterns is formally proven. It extends earlier works presented in [15,16,17], and asks how far can we understand and explain emergence following this track?

Keywords: sandpile models, fixed points, emergence.

1 Introduction

In the spirit of Chazelle’s natural algorithms [3], we propose a study of sandpile models via the tools of theoretical computer science. Sandpile configurations are words over the alphabet \mathbb{N} , where each letter is a number of stacked sand grains. The dynamics is described by local rules letting grains move from column to column. Historically, discrete sandpile models were first introduced by the physicists Bak, Tang and Wiesenfeld, as a paradigmatic example of *self-organized critical* systems in which instabilities are described by power laws [2]. The combinatorial study of one-dimensional models has been initiated by Goles in [9], and subsequently continued as a case where simple rules lead to surprisingly complex, yet tractable, mathematical problems on term rewriting systems. A particular focus has been given to the dynamics starting from a finite pile of sand grains (which is equivalent to adding them one by one, as they would fall in an hourglass), and to the structure of the language \mathcal{L} made of fixed point configurations reached at the end of the stabilization process [5,6,7,11].

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The model. Let $\mathbb{Z}_f^{\mathbb{N}}$ denote the set of integer sequence which are ultimately null (a sequence $h = (h_0, h_1, h_2, \dots)$ is ultimately null if there exists a positive integer i such that, for $j \geq i$, we have $h_j = 0$). We note by Δ the morphism $\mathbb{Z}_f^{\mathbb{N}} \rightarrow \mathbb{Z}_f^{\mathbb{N}}$ such that for any $h \in \mathbb{Z}_f^{\mathbb{N}}$ and any $i \in \mathbb{N}$, $\Delta h_i = h_i - h_{i+1}$. Notice that Δ is bijective. We say that h_i is the *number of grains* and Δh_i is *slope* of h in *column* i . The set \mathcal{C} is the subset of sequences of $\mathbb{Z}_f^{\mathbb{N}}$ which are non increasing.

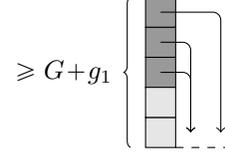


Fig. 1: The rule (2, 1) can be applied if and only if $\Delta h_i \geq 5$. For convenience, *right* is the direction of grains fall.

Definition 1. A decreasing sandpile model is a DDS defined by:

- **Configurations.** The set of configurations is the subset \mathcal{C} of $\mathbb{Z}_f^{\mathbb{N}}$,
- **Transition rule.** A p -tuple $\mathcal{R} = (g_1, g_2, \dots, g_p)$ with $g_1 \geq \dots \geq g_p > 0$. From a configuration h , a transition at i leads to the configuration h' , denoted $h \xrightarrow{i} h'$ when:
 - $h'_i = h_i - G$ with $G = \sum_{j=1}^p g_j$;
 - $h'_{i+j} = h_{i+j} + g_j$ for $1 \leq j \leq p$.

To lighten them, almost all notations do not mention the rule $\mathcal{R} = (g_1, \dots, g_p)$, as it does not change during the evolution. Since configurations must be non-increasing, the rule can be applied at column i if and only if $\Delta h_i \geq G + g_1$ (see Figure 1). Note that the rule application at i (which we also call *fire* at i) conserves the total number of sand grains. We also denote $h \rightarrow h'$ when the fired column is not given, and \rightarrow^* its reflexo-transitive closure. When $\Delta h_i \geq G + g_1$, we say that column i is *unstable*. A configuration is *stable*, or a *fixed point*, if it has no unstable column. Let \mathcal{C}_s be the set of stable configurations.

$$\mathcal{C}_s = \{h \in \mathcal{C} \mid \forall i \in \mathbb{N} : 0 \leq \Delta h_i < G + g_1\}.$$

We denote 0^ω the infinite sequence of 0. The models are non-deterministic: the rule is applied once at each time step. An example of successive applications for the rule (2, 1), ending in a fixed point, is given below,:

$$h = (26, 0^\omega) \xrightarrow{0} (23, 2, 1, 0^\omega) \xrightarrow{0} (20, 4, 2, 0^\omega) \xrightarrow{0} (17, 6, 3, 0^\omega) \xrightarrow{0} (14, 8, 4, 0^\omega) \\ \xrightarrow{0} (11, 10, 5, 0^\omega) \xrightarrow{1} (11, 7, 7, 1, 0^\omega) \xrightarrow{2} (11, 7, 4, 3, 1, 0^\omega).$$

Since Δ is a bijection on $\mathbb{Z}_f^{\mathbb{N}}$, a configuration h can be encoded by Δh , or $\Delta^2 h$. The sequence Δh of slopes is a more local (spatially uniform) representation of configurations. for the fixed point of the example above, $\Delta h = (4, 3, 1, 2, 1, 0^\omega)$. When not specified, the sequence of slopes is the default manner to give a configuration throughout the paper.

Let $\mathcal{C}(h) = \{h' \mid h \rightarrow^* h'\}$ denote the set of configurations *reachable* from h . We focus on the case when $h = (N, 0^\omega)$ *i.e.* configurations with a finite number $N \in \mathbb{N}$ of grains stacked on column 0, and no grains elsewhere. The set $\mathcal{C}((N, 0^\omega))$ is simply denoted by $\mathcal{C}(N)$.

The following theorem is obtained from *diamond property* plus *termination* (see for example [1] and [12]).

Theorem 1. $\mathcal{C}(h)$ endowed with \rightarrow has a graded lattice structure. Moreover, for any column i and any $h' \in \mathcal{C}(h)$, the number of firings of i during a sequence of transitions from h to h' does not depend of the chosen sequence of transitions.

In particular, even though decreasing sandpile models are non-deterministic, a unique fixed point, denoted by $\pi(h)$, is reached from any finite configuration h . The goal of this study is to understand the structure of $\pi(N) = \pi((N, 0^\omega))$. Configurations h' of $\mathcal{C}(h)$ admit another representation called *shot sequence*. It is the sequence $v = (v_i)_{i \in \mathbb{N}}$ where v_i is the number of times column i has been fired, from h to h' . The shot sequence of the stable configuration on our example is $v = (5, 1, 1, 0^\omega)$. Of course Δv or $\Delta^2 v$ can also encode a configuration of $\mathcal{C}(c)$. Surprisingly, the encoding $\Delta^2 v$ plays a fundamental role in the study of fixed points.

Remark 1. The classical one-dimensional sandpile rule is the 1-tuple (1), and the Kadanoff sandpile rule with parameter p is the p -tuple $(1, 1, \dots, 1)$. Important simplifications arise when the rule is a 1-tuple (g_1) , and a result similar to [9,10] is obtained with the same technics (by grouping grains in little stacks of g_1 units). Consequently, in the present work, we focus on the case $p > 1$.

The contribution. For a given decreasing sandpile rule \mathcal{R} , we are interested in the language (of sequences of slopes) of fixed points reached from initial configurations of the form $(N, 0^\omega)$.

$$\mathcal{L}_{\mathcal{R}} = \{\pi(N), N \in \mathbb{N}\} = \left(\bigcup_{N \in \mathbb{N}} \mathcal{C}(N) \right) \cap \mathcal{C}_s.$$

Experiments suggest that words from these languages all present regular repetitions of short patterns. Interestingly, when $p > 1$ those regular repetitions do not cover the entire non null part of the fixed point, but emerge from the dynamics: all the grains are initial stacked on column 0, and their toppling towards a stable configuration lets the regularities appear only on the right of some relatively small but asymptotically infinite position. Figure 2 presents some computer simulations, which lead to the following conjecture.

Conjecture 1. For any rule $\mathcal{R} = (g_1, \dots, g_p)$, there exists words $w_l, w_c, w_{c'}, w_r, w_f$ and $w_{f'}$, each of length at most p , such that if $h = \pi(N)$ then there exists an integer n in $\mathcal{O}(\log(N))$ such that $(\Delta h_i)_{i \geq n} = w_l^* (w_c + w_{c'} + \epsilon) w_r^* (w_f + w_{f'}) 0^\omega$, where $*$ denotes finite repetitions, $+$ the *or*, and ϵ the empty word.

There also exists words u_l, u_r on the alphabet $\{0, 1\}$, each of length at most p , such that if v is the shot sequence of $\pi(N)$, then there exists an integer n' in $\mathcal{O}(\log(N))$ such that $(\Delta^2 v_i)_{i \geq n'} = u_l^* (0 + \epsilon) u_r^* 0^\omega$.

Remark that the length $\mathcal{O}(\log N)$ of the irregular left part is negligible compared to the length $\Theta(\sqrt{N})$ of the non null part of the sequences (given a decreasing rule, it is not hard to notice that this part is in $\Theta(\sqrt{N})$ because the

Rule (4,1,1) and $\pi(4000)$:

$$\text{Delta } h = (\begin{array}{cccccccccccccccc} 8, & 6, & 6, & 2, & 8, & 9, & 7, & 9, & 5, & \underline{9}, & \underline{9} \end{array})$$

$$\text{Delta}^2 v = (\begin{array}{cccccccccccccccc} 188, & 84, & -59, & 7, & 9, & -3, & 1, & 1, & 1, & 1, & 1, & 1, & 1, & 1, & 1, & 1, & 1, & 1, & 1, & 0^{\omega} \end{array})$$

Rule (3,2,1,1) and $\pi(3800)$:

$$\text{Delta } h = (\begin{array}{cccccccccccccccccccc} 3, & 8, & 2, & 6, & 6, & 9, & 3, & 8, & 9, & 9, & 8, & 9, & 7, & 2, & 8, & 4, & \underline{9}, & 5, & \end{array})$$

$$\text{Delta}^2 v = (\begin{array}{cccccccccccccccc} 332, & -24, & 39, & -62, & 29, & -4, & 4, & -4, & 3, & 0, & 1, & 0, & 0, & 1, & 0, & 1, & 0, & 1, & 0, & 1, & 0^{\omega} \end{array})$$

Fig. 2: Examples of fixed points with the regularly repeated pattern highlighted.

fixed point is a rectangular triangle of area N and bounded slope (by stability from above and from below since there are no plateau of length greater than $p + 1$). Thus the conjecture gives a nearly exhaustive characterization of fixed points. Also remark that the first part of the conjecture is a (spectacular but direct) consequence of the second part. Thus, in the study, we will focus in the second part.

For the Kadanoff sandpile models, this conjecture has been proven in [15,16,17]. In the present paper we generalize in a large way the result about fixed points on Kadanoff sandpiles (note that Kadanoff sandpiles already generalizes the classical one-dimensional model of [9,10]). We treat the extreme cases: when g_1 is sufficiently large to command the evolution (Theorem 4), and when g_1 is sufficiently small (Theorem 5).

First, the language $\mathcal{L}_{\mathcal{R}}$ is investigated via arguments of linear algebra linked to static constraints obtained from the definition of a fixed point (Section 2). By static constraints, we mean constraints telling that the configuration is stable, but nothing (or very few) about its creation process. This step gives an automaton whose set of accepting words contains $\mathcal{L}_{\mathcal{R}}$. Then, we use a dynamic and recursive construction of fixed points to refine the result. More precisely, the fixed point $\pi(N + 1)$ is obtained from $\pi(N)$ by adding a single grain on column 0 and performing all the possible firings until stability is reached (this process is called the $(N + 1)^{\text{th}}$ avalanche). Using combinatorial arguments, we formally prove the emergence of regular patterns on fixed points, when the rule verifies arithmetic properties (Section 3).

Notice that those two steps informally correspond to the hybrid nature of sand. The algebraic part considers it as a (quasi) continuous fluid, while the combinatorial one considers it as a set of discrete atoms. To achieve this generalization, we made major improvements of the developments around the Kadanoff model. The first one is the use of stronger algebraic theorems about matrix norms, to ensure a convergence of the algebraic part. The second one is the introduction of the encoding $\Delta^2 v$ of configurations, which is very useful to describe fixed points and avalanches, and allows to easily consider these two notions simultaneously. For the Kadanoff model, the automaton is quite simple (and so is the second step), thus the tools we introduce in the present work were not explicitly defined and used previously [15,16,17].

Finally, we discuss the results and conjecture their generalization to all decreasing sandpile models (Section 4). We strongly conjecture that our approach can succeed for any decreasing sandpile model. But, currently, we do not have a general proof, even if we are able to prove the conjectures for numerous particular rules.

2 Static study of the fixed points

Equivalent representations. Decreasing sandpile models can also be seen as *chip-firing games* (basically with a vertex for each column and a chip content corresponding to the slope) with a sink, what Dhar calls *abelian sandpile models* in [4]. The same notion of *equivalence class* applies. We define it on the set $\mathbb{Z}_f^{\mathbb{N}}$ of all ultimately null configurations. For all $h, h' \in \mathbb{Z}_f^{\mathbb{N}}$, we have $h \equiv_{\mathcal{R}} h'$ if and only if h' can be reached from h by a sequence of *firings* and *anti-firings* (the reverse of a firing) on column indices in \mathbb{N} (configurations may have columns with negative sand content). It is reflexive, symmetric and transitive, hence an equivalence relation. The equivalence class of a configuration h is denoted $[h]$. Let

$$\Pi(h) = [h] \cap \mathcal{C}_s \quad \text{and} \quad \mathcal{L}'_{\mathcal{R}} = \bigcup_{N \in \mathbb{N}} \Pi(N).$$

The developments of this section will give clues on the language $\mathcal{L}'_{\mathcal{R}}$. Note that $\pi(N) \in \Pi(N)$, hence $\mathcal{L}_{\mathcal{R}} \subseteq \mathcal{L}'_{\mathcal{R}}$, therefore the conclusions of Section 2 will apply to our main language of interest. The arguments are based on static equations, talking about the final stability of the configurations of $\Pi(N)$. The main static relation we exploit is simple to notice. Let us canonically extend the definition of shot sequence to the configurations of $[N]$. This representation is obviously linked to the sequence of heights, by the equality below, which simply expresses the grain balance on column i after firings and anti-firings, in a decreasing sandpile rule $\mathcal{R} = (g_1, \dots, g_p)$ with $p > 1$.

$$\text{For all } i \geq p \text{ we have } h_i = -G v_i + g_1 v_{i-1} + g_2 v_{i-2} + \dots + g_p v_{i-p}. \quad (1)$$

In this section, we will basically make use of Equation (1) and the constraint that the configurations in $\Pi(N)$ are stable, non-increasing, and ultimately null. From this we will first construct a recurrence equation describing the fixed point from left to right: given $(\Delta v_j)_{i \leq j < i+p}$ we will express $(\Delta v_j)_{i+1 \leq j < i+p+1}$. Then arguments of linear algebra will be employed in order to prove a convergence result on iterations of this system: columns on the right of a position in $\mathcal{O}(\log N)$ have strong regularity properties. The results will eventually be expressed as an automaton (depending on \mathcal{R}) recognizing the sequence of slopes of any configuration in $\Pi(N)$. Talking about the language $\mathcal{L}'_{\mathcal{R}}$ is not a goal in itself, but a way to express the fact the the arguments we exploit in this section are quite general: using static arguments we set up restrictions on the language of sequences of slopes for any configuration in $\Pi(N)$.

Section 3 will present a different and complementary kind of arguments, linked to the dynamic of the model, that will allow to refine the results to

get a precise characterization of $\mathcal{L}_{\mathcal{R}}$ when the rule \mathcal{R} verifies some arithmetic properties.

The perturbed weighted mean system. Equation (1) can be expressed in terms of Δh and Δv , leading to the *recurrence relation*

$$\Delta v_i = \frac{1}{G} \left(g_1 \Delta v_{i-1} + g_2 \Delta v_{i-2} + \cdots + g_p \Delta v_{i-p} \right) - \frac{\Delta h_i}{G}. \quad (2)$$

In matricidal form, it gives a system from \mathbb{Z}^p to \mathbb{Z}^p that we call *perturbed weighted mean system*:

$$\Delta V_i = M \Delta V_{i-1} - \frac{\Delta h_i}{G} K \quad (3)$$

with

$$\Delta V_i = \begin{pmatrix} \Delta v_{i-p+1} \\ \vdots \\ \Delta v_{i-1} \\ \Delta v_i \end{pmatrix} \quad M = \begin{pmatrix} 0 & 1 & & 0 \\ & & \ddots & \\ 0 & 0 & & 1 \\ \frac{g_p}{G} & \frac{g_{p-1}}{G} & \cdots & \frac{g_1}{G} \end{pmatrix} \quad K = \begin{pmatrix} 0 \\ \vdots \\ 0 \\ 1 \end{pmatrix}.$$

It is composed of two parts. First a linear map $M : \mathbb{R}^p \rightarrow \mathbb{R}^p$ that

- shifts all the values one row upward;
- for the last component, computes the mean of ΔV_{i-1} weighted by (g_1, \dots, g_p) .

Second a discrete perturbation $\frac{\Delta h_i}{G}$ subtracted to the last component, so that $\Delta v_i \in \mathbb{Z}$. Intuitively, iterating the system, *i.e.*, computing the weighted mean and subtracting a bounded perturbation $((\Delta h_i)_{i \in \mathbb{N}} \in \mathcal{C}_s)$ will quickly tend to output values which are close to each other. In other words, the sequence $(\Delta V_i)_{i \geq p-1}$ will tend to uniform vectors. This is what we are about to prove. For convenience, let $m_i = \frac{1}{G}(g_p, \dots, g_1) \Delta V_i$ denote the weighted mean of ΔV_i , and \underline{m}_i (resp. \bar{m}_i) the minimal (resp. maximal) value of ΔV_i .

As a starting point of the system, we define $\Delta V_{-1} = {}^t(\frac{N}{g_p}, 0, \dots, 0, -v_0)$, emulating the fact that column 0 receives N units of sand grains. Now Equation (2) and the perturbed weighted mean system hold for all $i \in \mathbb{N}$. For the configurations of $\Pi(N)$ we have

$$\frac{N}{G + g_1} \leq v_0 \leq \frac{N}{G}, \quad (4)$$

because from the initial configuration we have to perform at least the left bound number of firings (less and the configuration cannot be stable), and at most the right bound number of firings (more and it is not possible to get a non-increasing and ultimately null configuration).

Regarding the discrete perturbation $\frac{\Delta h_i}{G}$, for all configurations of the set $\Pi(N)$ we have a relation for stability (left) and a relation for integrity (right).

$$0 \leq \frac{\Delta h_i}{G} < 1 + \frac{g_1}{G} \leq 2 \quad \text{and} \quad \Delta h_{i+1} \equiv G m_i \pmod{G}. \quad (5)$$

Note that consequently, for a given m_i there are at most two possible values of Δh_i that match these two constraints, and only one when $m_i \geq G + g_1$. This will be a key point in the construction of automata at the end of this section.

For the convergence of $(\Delta V_i)_{i \in \mathbb{N}}$ towards uniform vectors, we first prove that it converges in $\mathcal{O}(\log N)$ iterations to vectors of bounded amplitude.

Lemma 1. *There exists a constant α and a $n_0 \in \mathcal{O}(\log N)$ s.t. $\bar{m}_i - \underline{m}_i < \alpha$ for all $i > n_0$.*

Proof (sketch). Equation (4) implies that $\bar{m}_{-1} - \underline{m}_{-1}$ is in $\Theta(N)$. In order to prove that iterations of the perturbed weighted mean system tend to uniform vector, that is, each value will become closer and closer to the mean value, we take $M_i = {}^t(m_i, \dots, m_i) \in \mathbb{R}^p$ and study the sequence $(Z_i)_{i \in \mathbb{N}}$ where $Z_i = \Delta V_i - M_i$, which converges to $\mathbb{0}_p = {}^t(0, \dots, 0) \in \mathbb{Z}^p$. From Equation (3) we get a relation of the form

$$Z_i = O Z_{i-1} - \frac{\Delta h_i}{G} L$$

where O is a contracting map: its spectral radius is strictly smaller than 1 (proved with a classical result due to Eneström and Kakeya, see for example [8]). As a consequence, we can isolate the contracting map and the sum of perturbations,

$$Z_n = O^{n+1} Z_{-1} + \frac{1}{G} \sum_{i=0}^n \Delta h_i O^{n-i} L.$$

The left part of the sum tends exponentially to $\mathbb{0}_p$ (see for example [13] for a discussion on contracting maps), and the right part is upper bounded by some constant $\alpha - 1$. Since the norm of Z_i is in the order of $\bar{m}_i - \underline{m}_i$, the norm of Z_{-1} is in $\Theta(N)$ and there exists an iteration n_0 in $\mathcal{O}(\log N)$ such that the left term is strictly smaller than 1, leading to the result. \square

Secondly, it converges, at least linearly in the amplitude, to a sequence which is non-increasing and where two consecutive values are equal or differ by one.

Lemma 2. *There exists d in $\mathcal{O}(\bar{m}_i - \underline{m}_i)$, such that for all $k \geq i + d$, we have $\Delta v_k - \Delta v_{k+1} \in \{0, 1\}$. Moreover, $\Delta v_{k+1} = \Delta v_k - 1$ only if $m_i - \underline{m}_i < \frac{g_1}{G}$.*

Proof (sketch). The core argument of this proof is that, when $\bar{m}_i \neq \underline{m}_i$, the weighted mean m_i is strictly between \bar{m}_i and \underline{m}_i . The perturbation (bounded by Equation (5)) subtracted to it then always leads to an integer Δv_{i+1} which is strictly below \bar{m}_i and greater or equal to $\underline{m}_i - 1$. As a consequence, the result holds if the sequence $(\Delta v_k)_{k \geq i+d}$ is non-increasing. In order to prove this, we can first notice that while Δv_{i+1} is above or equal to \underline{m}_i , it is still strictly below \bar{m}_i so ΔV_i tend linearly to uniform vectors. When it happens that $\Delta v_{i+1} = \underline{m}_i - 1$, the values embedded in ΔV_i must already be very close to each other, in order for the mean m_i to be just a little bit above \underline{m}_i , so that the perturbation (strictly smaller than 2 from Equation (5)) subtracted to it can lead to a Δv_{i+1} strictly below \underline{m}_i . In this case, a careful look at the constraints on the closeness of the values embedded in ΔV_i allows to conclude. \square

The combination of Lemmas 1 and 2 gives the expected result, which we will express in term of the second derivative of the shot sequence. Let $\Delta^2 V_i = {}^t(\Delta^2 v_{i-p+1}, \dots, \Delta^2 v_{i-1})$. Note that $\Delta^2 V_i$ can be computed from ΔV_i , and so can the difference $m_i - \underline{m}_i$, with the function m defined as

$$m(q_1, \dots, q_{p-1}) = \sum_{1 \leq j < k \leq p} q_{p-j} g_k = \sum_{j=1}^{p-1} \left[q_{p-j} \left(\sum_{k=j+1}^p g_k \right) \right].$$

Property 1. – For the product order, if $U < U'$ then $m(U) < m(U')$.

– If $\Delta v_i = \underline{m}_i$ then $m(\Delta^2 V_i) = G(m_i - \underline{m}_i)$.

– $m(0, 0, \dots, 0, 0, 1) = g_2 + g_3 + \dots + g_p$; $m(0, 0, \dots, 0, 1, 0) = g_3 + \dots + g_p$;

$m(1, 0, \dots, 0, 0, 0) = g_p$; $m(1, 1, \dots, 1, 1, 1) = g_2 + 2g_3 + \dots + (p-1)g_p$.

Proposition 1. *There exists a column n_1 in $\mathcal{O}(\log N)$, such that*

$$\text{for all } i \geq n_1 \text{ we have } \Delta^2 v_i \in \{0, 1\}.$$

Moreover, $\Delta^2 v_i = 1$ only if $m(\Delta^2 V_i) < g_1$, and

$$\text{for all } i \geq n_1 \text{ we have } \Delta h_{i+1} = m(\Delta^2 V_i) + \Delta^2 v_i G.$$

Proof. First part is a straight combination of Lemmas 1 and 2. For the middle part, when Δv is non-increasing we always have $\Delta v_i = \underline{m}_i$ and Property 1 applies. For the last part, from Relation (3) we have $\Delta h_{i+1} = G(m_i - \Delta v_{i+1})$. When $i \geq n_1$ the first part states that $\Delta^2 v_i$ equals 0 or 1 (recall that $\underline{m}_i = \Delta v_i$), and in both cases the conclusion is reached by applying Property 1. \square

Remark 2. To avoid confusion, remark that $\Delta^2 v_i$ is not the last component of $\Delta^2 V_i$, but the last component of $\Delta^2 V_{i+1}$.

The automaton. Proposition 1 gives restrictions on the language $\mathcal{L}'_{\mathcal{R}}$. Let us now express these restrictions in the framework of automata theory. Given a decreasing sandpile rule, its *recurrence automaton* will be a Muller automaton (a kind of Büchi automaton with stronger accepting condition, see [14] for a definition) recognizing a subset of \mathcal{C}_s which includes $\mathcal{L}'_{\mathcal{R}}$ (asymptotically).

States of the automata correspond to vectors $\Delta^2 V_i \in \{0, 1\}^{p-1}$, and transitions to iterations of the perturbed weighted mean system. Depending on the value of $m(\Delta^2 V_i)$, Proposition 1 tells that there is one or two out-going transitions from a state. We label a transition with the value of Δh_{i+1} given by Proposition 1. We consider the whole set of states as potential initial states, and 0^{p-1} as the unique entry in the acceptance table of Muller automata, *i.e.*, runs must end in an infinite loop on the state 0^{p-1} , corresponding to the fact that all words of $\mathcal{L}'_{\mathcal{R}}$ are ultimately null (by definition).

Definition 2. *Given a decreasing sandpile rule \mathcal{R} , let $\mathcal{A}_{\mathcal{R}}$ be its recurrence automaton, which is the Muller automaton whose set of states is $\{0, 1\}^{p-1}$, alphabet is \mathbb{N} , set of initial states is $\{0, 1\}^{p-1}$, acceptance table is $\langle \{0^{p-1}\} \rangle$, and there exists a transition $(q_1, \dots, q_{p-1}) \xrightarrow{a}_{\mathcal{R}} (q'_1, \dots, q'_{p-1})$ if and only if*

- $(q'_1, \dots, q'_{p-2}) = (q_2, \dots, q_{p-1})$;
- either $q'_{p-1} = 0$
or $q'_{p-1} = 1$ and $m(q) < g_1$;
- $a = m(q) + q'_{p-1}G$.

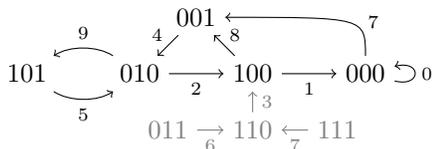


Fig. 3: Recurrence automata $\mathcal{A}_{\mathcal{R}}$ for $\mathcal{R} = (3, 2, 1, 1)$, where irrelevant states are shaded. Note that it is in accordance with the example from Figure 2.

An example of recurrence automaton is given on Figure 3. Theorem 2 rephrases Proposition 1 in order to characterize $\mathcal{L}'_{\mathcal{R}}$.

Theorem 2. *Let $\mathcal{L}(\mathcal{A}_{\mathcal{R}})$ denotes the language of infinite words recognized by $\mathcal{A}_{\mathcal{R}}$. For all $\Delta h \in \Pi(N)$, there exists a column n_2 in $\mathcal{O}(\log N)$ such that*

$$(\Delta h_i)_{i \geq n_2} \in \mathcal{L}(\mathcal{A}_{\mathcal{R}}).$$

Theorem 2 is equally valid if we simplify recurrence automata to their *ultimately relevant* states, *i.e.*, those belonging to a directed cycle (note that 0^{p-1} is always an ultimately relevant state with a loop labelled with slope value 0). Let $\mathcal{A}'_{\mathcal{R}}$ denote these simplified automata. Languages $\mathcal{L}(\mathcal{A}'_{\mathcal{R}})$ contain suffixes of the languages we are interested in. It is not tight enough to fit $\mathcal{L}_{\mathcal{R}}$, not even exactly $\mathcal{L}'_{\mathcal{R}}$, but can nevertheless be considered as an important progress compared to \mathcal{C}_s . Furthermore, the developments so far are only based on *static* relations coming from the model definition. Finally, let us recall that $w(\pi(N)) \in \Theta(\sqrt{N})$ for all rule and number of grains N , therefore $\mathcal{L}(\mathcal{A}'_{\mathcal{R}})$ contains prefixes of $\mathcal{L}'_{\mathcal{R}}$ that asymptotically account for the whole fixed points. Note that this is not an equality: $\mathcal{A}'_{\mathcal{R}}$ recognizes other words.

3 Dynamic study of the fixed points

Avalanches. This section complements the developments around $\mathcal{L}_{\mathcal{R}}$ presented so far ($p > 1$), with arguments linked the dynamics of decreasing sandpile models. It is based on the fact that for a given rule, $\pi(0), \pi(1), \pi(2), \dots, \pi(N)$ must all belong to $\mathcal{L}(\mathcal{A}'_{\mathcal{R}})$ (starting from some index n_2 in $\mathcal{O}(\log N)$). Let us present the notion of *avalanche*, and two cases where combinatorial arguments allow to characterize asymptotically $(\Delta^2 v_i)_{i \in \mathbb{N}}$ (hence $(\Delta h_i)_{i \in \mathbb{N}}$, see Section 4) when g_1 is above some upper threshold or below some lower threshold (note that the second case includes Kadanoff sandpile models).

Let $c^{\downarrow 0}$ denote the configuration obtained from c by adding one grain on column 0. If $c \xrightarrow{i_1} \dots \xrightarrow{i_k} c'$ then $c^{\downarrow 0} \xrightarrow{i_1} \dots \xrightarrow{i_k} c'^{\downarrow 0}$. This is in particular true for $c = (N, 0^\omega)$ and $c' = \pi(N)$, thus fixed points can inductively be computed with

$$\pi(N+1) = \pi(\pi(N)^{\downarrow 0}),$$

by, from $\pi(0) = (0^\omega)$, repeating N times the addition of one grain followed by the stabilization process. We will now use this point of view. Let the $(N+1)^{th}$ avalanche be the sequence $(a_i)_{i \geq 0}$ where a_i denotes the number of times that column i has been fired in the stabilization process from $\pi(N)^{10}$ to $\pi(N+1)$. Let plain h, v (resp. primed h', v') denote representations of $\pi(N)$ (resp. $\pi(N+1)$),

$$(v'_i - v_i)_{i \geq 0} = (a_i)_{i \geq 0} \quad \text{and} \quad (\Delta^2 v'_i - \Delta^2 v_i)_{i \geq 0} = (\Delta^2 a_i)_{i \geq 0}. \quad (6)$$

Avalanches are studied in [15], which is generalized as follows.

Theorem 3. 1. For each integer $i \geq 1$, we have $a_i \in \{0, 1\}$.

2. If $a_j = a_{j+1} = \dots = a_{j+p-1} = 0$ for some $j \in \mathbb{N}$, then $a_k = 0$ for all $k \geq j$.

3. If $\Delta h_{j+r} + g_{r+1} < G + g_1$ for some $j \in \mathbb{N}$ and all $0 \leq r \leq p-1$, then integer j satisfies condition 2 above.

Proof (sketch). Arguments similar to [15], telescoping sum for third point. \square

First, note that Proposition 1 and Equation (6) imply that there exists $n_2 \in \mathcal{O}(\log N)$ such that $(\Delta^2 a_i)_{i \geq n_2} \in \{-1, 0, 1\}^\omega$, hence there is no factor 010 or 101 in $(a_i)_{i \geq n_2}$. A simple case by case study of what does the value of $\Delta^2 a_i$ implies on a_i, a_{i+1} and a_{i+2} shows that, up to a one unit shift, the patterns $1(-1)$ and $(-1)1$ of $(\Delta^2 a_i)_{i \geq n_2}$ are delimiters of the intervals of 0 and 1 in $(a_i)_{i \geq n_2}$. Moreover, the length of intervals of 0 in $(a_i)_{i \geq n_2}$ is at most $p-1$ (except the ultimate 0^ω). Basic considerations of this type lead to the following proposition (* is the Kleene star denoting finite repetitions).

Proposition 2. $(\Delta^2 a_i)_{i \geq n_2}$ is suffix of a sequence in $(0^* 1(-1) 0^* (-1) 1)^* 0^\omega$.

We will have two main arguments: the stopping condition of Theorem 3, and the compatibility between $\pi(N)$ and $\pi(N+1)$ (Proposition 1 and Theorem 2). For $1 \leq i \leq p-1$, let $E_i \in \{0, 1\}^{p-1}$ have all null components except the i^{th} .

Case $g_1 > m(\mathbf{1}, \mathbf{1}, \dots, \mathbf{1})$. In this case two transitions are possible from any state in the automaton *i.e.*, Theorem 2 gives no constraint other than $\Delta^2 v_i \in \{0, 1\}$ for all $i \geq n_2$. Nevertheless we show that, on the right of column n_2 , the avalanche process fires a set of consecutive columns and stops.

Lemma 3. If $\exists k \in \mathbb{N}$ such that $\Delta^2 v_{n_2+k} = 0$, then $a_j = 0$ for $j \geq n_2 + k + 1$.

Proof (sketch). The goal is to reach the stopping condition of Theorem 3. From the hypothesis, $\Delta h_{n_2+k+1} = m(\Delta^2 V_{n_2+k}) + \Delta^2 v_{n_2+k} G \leq m(\mathbf{1}, \dots, \mathbf{1}) < g_1$. For $2 \leq r \leq p$ we have $\Delta^2 V_{n_2+k+r-1} \leq {}^t(1, \dots, 1) - E_{p+1-r}$ and $m(E_{p+1-r}) = g_r + g_{r+1} + \dots + g_p$. Since m is a linear map, $\Delta h_{n_2+k+r} + g_r \leq m(\mathbf{1}, \dots, \mathbf{1}) + g_1 + g_2 + \dots + g_r < g_1 + G$ and from Theorem 3 the avalanche stops (item 3). \square

Theorem 4. $(a_i)_{i \geq n_2} \in 1^* 0^\omega$ and $(\Delta^2 v_i)_{i \geq n_2} \in 1^* (0 + \epsilon) 1^* 0^\omega$.

Proof (sketch). The result is proven by induction on N , the base case is obvious. For the induction, we have $(\Delta^2 v_i)_{i \geq n_2} \in 1^{k'} (0 + \epsilon) 1^{k'} 0^\omega$. From Proposition 2 and Equation 6, we have either $(\Delta^2 a_i)_{i \geq n_2} = 0^\omega$, or $(\Delta^2 a_i)_{i \geq n_2} = 0^{k-1} (-1) 1 0^\omega$, otherwise it contradicts Proposition 1. In both cases the statement holds. \square

Case $m(E_r) \leq g_1 < m(E_{r+1})$ with $r + 1 \leq \frac{p}{2}$. Note that it is always true that $m(E_1) = g_p \leq g_1$. In this case, we will prove two preliminary lemmas and a description of $(\Delta^2 v_i)_{i \geq n_2}$ similar to Theorem 4.

Lemma 4. *In $(\Delta^2 v_i)_{i \geq n_2}$ two 1 are separated by at least $p - r - 1$ values 0.*

Proof (sketch). In terms of the automaton $\mathcal{A}_{\mathcal{R}}$, while the state q ends with strictly less than $p - r - 1$ values 0, there is only one transition (adding a new 0 at the end of q) because $m(q) \geq g_1$ (the minimal case for m is $q = 0^r 1 0^{p-r-2} = E_{r+1}$). The result follows by induction. \square

The next Lemma tells that $p - r$ values 0 in $\Delta^2 v$ stop the avalanche.

Lemma 5. *If $\Delta^2 v_j = 1$ and $\Delta^2 v_{j+p-r} = 0$ for $j > n_2 + p$, then $(a_i)_{i \geq j+2} = 0^\omega$.*

Proof (sketch). The goal is again to reach the stopping condition of Theorem 3. From Lemma 4, the hypothesis, and because $m(E_{p-1}) + g_1 = G$, we have $\Delta h_{j+2} + g_1 = m(\Delta^2 V_{j+1}) + g_1 \leq m(E_{p-1}) + m(E_{r-1}) + g_1 < G + g_1$. It follows from similar arguments that $\Delta h_{j+k} + g_{k-1} < G + g_{k-1}$ for $2 \leq k \leq p + 1$, thus Theorem 3 (item 3) applies. \square

Theorem 5. *$(a_i)_{i \geq n_2+p} \in 1^* 0^\omega$ and $(\Delta^2 v_i)_{i > n_2+p}$ is suffix of an element in $(0^{p-r-1} 1)^* (0 + \epsilon) (0^{p-r-1} 1)^* 0^\omega$.*

Proof (sketch). The result is proven by induction on N . The base case is obvious, and the induction is obtained from Equation (6), and the constraints given by Proposition 1, Lemma 4 and Lemma 5. \square

4 Conclusions and perspectives

This paper has explored the languages of sequence of slopes for fixed points of decreasing sandpile models, $\mathcal{L}_{\mathcal{R}}$. After a general development based on linear algebra, in two cases, Theorems 4 and 5 give precise asymptotic characterizations of $\Delta^2 v$, which apply to the sequences of slopes (via Proposition 1). We conjecture that a similar characterization holds for any decreasing sandpile rule (see Conjecture 1).

Note that for the two cases we solved, we have $w_l = w_r$ (from the labeling of recurrence automata and Theorem 2), but this is not always the case. For example, we strongly conjecture that for the rule $(6, 1, 1, 1, 1)$, the sequence $\Delta^2 v$ is asymptotically of the form $(0011)^{k_l} (0 + \epsilon) (01)^{k_r}$ and a subsequent avalanche leads to $(1100)^{k_l-1} 0 (01)^{k_r+2}$ when the central part was 0, and to $(1100)^{k_l+1} (01)^{k_r-2}$ when the central part was ϵ . Hence w_l and w_r may differ.

The proof technic works in several stages, that may highlight the fact that sandpiles are at the edge between discrete and continuous systems. Lemma 1 uses arguments of linear algebra corresponding to a rough continuous nature, while Lemma 2 and Section 3 refine the study with precise combinatorial arguments corresponding to discrete dynamical phenomena.

The structure of avalanches is also an interesting point of view. We propose the following conjecture, satisfied for the cases treated completely in Section 3.

Conjecture 2. For any decreasing sandpile rule \mathcal{R} and any positive integer N , there exists n in $\mathcal{O}(\log N)$ such that the N^{th} avalanche verifies $(a_i)_{i \geq n} \in 1^* 0^\omega$.

According to the conjectures above, fixed points and avalanches are characterized on the right of some column n in $\mathcal{O}(\log N)$ compared to their width in $\Theta(\sqrt{N})$ *i.e.*, asymptotically completely. Though its relative size tends to be null, the unknown part between 0 and n is not bounded. The conjectures mean that as we add grains one by one, they trigger avalanches that let grains create and maintain regular patterns, after a transitional phase of unbounded length. Let us finish on a question: should this process be called *self-organized emergence*?

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A Muller automata

Muller automata were introduced in [14]. They are very close to Büchi automata and make use of the concepts of *automata* and *runs*.

Definition 3. An automaton is a 4-tuple $(\mathcal{Q}, \Sigma, \rightarrow, \mathcal{S})$, where

- \mathcal{Q} is a set of states;
- Σ is an alphabet;
- \mathcal{S} is a set of initial states;
- $\rightarrow \subseteq \mathcal{Q} \times \Sigma \times \mathcal{Q}$ is a set of transitions, and $(q, a, q') \in \rightarrow$ is denoted $q \xrightarrow{a} q'$.

An automaton describes a dynamic that is captured by *runs*.

Definition 4. For an automaton $\mathcal{A} = (\mathcal{Q}, \Sigma, \rightarrow, \mathcal{S})$ and an infinite input word $w \in \Sigma^\omega$, a run of \mathcal{A} on w is an infinite sequence of states $\rho \in \mathcal{Q}^\omega$ starting at some $\rho_0 \in \mathcal{S}$ and such that for all $i \in \mathbb{N}$ we have $\rho_i \xrightarrow{w_i} \rho_{i+1}$.

The definition of Muller automata adds an accepting condition for the recognition of infinite words.

Definition 5. A Muller automaton is a pair $(\mathcal{A}, \mathcal{T})$, where $\mathcal{A} = (\mathcal{Q}, \Sigma, \delta, \mathcal{Q}_{in})$ is an automaton and $\mathcal{T} = \langle \mathcal{Q}_1, \mathcal{Q}_2, \dots, \mathcal{Q}_k \rangle$ is an acceptance table with $\mathcal{Q}_i \subseteq \mathcal{Q}$ for $i \in \{1, 2, \dots, k\}$.

The Muller automata accepts (or recognizes) an input $w \in \Sigma^\omega$ if there is a run ρ of \mathcal{A} on w such that $\{q \in \mathcal{Q} \mid \exists^\omega n : \rho_n = q\} = F_i$ for some i .

The accepting condition is quite strong: for a word to be accepted, there must exist a run ρ and an entry F_i of the acceptance table, such that the set of states visited infinitely often by ρ is exactly F_i .

B Complete proofs

Proof (Lemma 1). From Equation (4), we have that $\bar{m}_{-1} - \underline{m}_{-1}$ is in $\Theta(N)$.

We introduce the matrix

$$D = \begin{pmatrix} 1 & 0 \\ \cdot & \cdot \\ 0 & 1 \end{pmatrix} - \frac{1}{G} \begin{pmatrix} g_p & g_{p-1} & \dots & g_1 \\ \vdots & \vdots & & \vdots \\ g_p & g_{p-1} & \dots & g_1 \end{pmatrix}.$$

For each vector U of \mathbb{R}^p , DU is obtained from U by subtracting to each component the weighted average $\frac{1}{G}(g_p, \dots, g_1)U$. Hence D is the matrix of the projection on the hyperplane formed by vectors U such that $(g_p, \dots, g_1)U = 0$, according to the direction of the vector $\mathbb{1}_p$.

Let $Z_i = D \Delta V_i$ and $M_i = m_i \mathbb{1}_p$, i.e., M_i is the uniform vector with all its components equal to $m_i = \frac{1}{G}(g_p, \dots, g_1) \Delta V_i$. We have $M_i = \Delta V_i - Z_i$.

To prove the lemma, it suffices to show an n_0 in $\mathcal{O}(\log N)$ such that $\|Z_{n_0}\|_\infty \leq \frac{\alpha}{2}$ for some constant α . We have

$$Z_i = D \Delta V_i = D \left(M \Delta V_{i-1} - \frac{\Delta h_i}{G} K \right) = DM \Delta V_{i-1} - \frac{\Delta h_i}{G} DK.$$

Since $\Delta V_i = Z_i + M_i$, and $DMM_i = DM_i = \mathbb{0}_p$, we have the recurrence equation

$$Z_i = DM (Z_{i-1} + M_{i-1}) - \frac{\Delta h_i}{G} DK = DM Z_{i-1} - \frac{\Delta h_i}{G} DK,$$

from which we get the expression

$$Z_n = (DM)^{n+1} Z_{-1} + \frac{1}{G} \sum_{i=0}^n \Delta h_i (DM)^{n-i} DK.$$

There exists a matrix norm (see *e.g.* [E. Isaacson and H. Keller. *Analysis of Numerical Methods*. Wiley, 1966]) such that for each $x \in \mathbb{R}$, for each pair A, B of matrices, and each vector U , we have $\|A + B\| \leq \|A\| + \|B\|$, $\|xA\| = |x| \|A\|$, $\|AB\| \leq \|A\| \|B\|$, $\|AU\|_\infty \leq \|A\| \|U\|_\infty$.

We can deduce that

$$\|Z_n\|_\infty \leq \|DM\|^{n+1} \|Z_{-1}\|_\infty + \frac{1}{G} \sum_{i=0}^n \Delta h_i \|DM\|^{n-i} \|DK\|_\infty,$$

from which, since for each integer i it holds that $\Delta h_i \leq G + g_1$, we get

$$\|Z_n\|_\infty \leq \|DM\|^{n+1} \|Z_{-1}\|_\infty + \frac{G + g_1}{G} \left(\sum_{j=0}^n \|DM\|^j \right) \|DK\|_\infty.$$

We now have to study DM . We recall that, for each (complex) square matrix A , the *spectrum* of A , denoted by $Sp(A)$, denotes the set of eigenvalues of A , and the *spectral radius* of A , denoted by $\rho(A)$, is the maximum among the moduli of values of $Sp(A)$.

Fact 1 *For each eigenvalue λ of $Sp(M) \setminus \{1\}$, we have $\frac{1}{G} \leq |\lambda| \leq \frac{G-1}{G} < 1$. Moreover, we have*

$$Sp(DM) \setminus \{0\} \subseteq Sp(M) \setminus \{1\}.$$

Let us postpone for a moment the demonstration of the fact above. The fact is used to obtain, as a direct corollary, that

$$\rho(DM) < 1.$$

For any $\epsilon > 0$ and any matrix A , the matrix norm satisfying the conditions above can be chosen in such a way that $\|A\| < \rho(A) + \epsilon$ (see *e.g.* [E. Isaacson and H. Keller. *Analysis of Numerical Methods*. Wiley, 1966] Th. 3, p. 12). Applied with $A = DM$ and ϵ such that $\rho(DM) + \epsilon < 1$, we get $\|DM\| < 1$. Thus

$$\|Z_n\|_\infty \leq \|DM\|^{n+1} \|Z_{-1}\|_\infty + \frac{G + g_1}{G(1 - \|DM\|)} \|DK\|_\infty.$$

Since $\|DM\| < 1$ and $\|Z_{-1}\|_\infty$ is in $\mathcal{O}(N)$, there exists n_0 in $\mathcal{O}(\log(N))$ such that $\|DM\|^{n+1} \|Z_{-1}\|_\infty \leq 1$, therefore we get the result with $\frac{\alpha}{2} = 1 + \frac{G+g_1}{G(1-\|DM\|)} \|DK\|_\infty$. \square

Proof (Fact 1). In order to prove Fact 1, let us first concentrate on the matrix M , which is a companion matrix of characteristic polynomial

$$\begin{aligned} x^p - \frac{1}{G} \sum_{k=0}^{p-1} g_{p-k} x^k &= (x-1) R(x) \\ \text{with } R(x) &= \frac{1}{G} \sum_{k=0}^{p-1} \left(\sum_{j=0}^k g_{p-j} \right) x^k \\ &= \frac{1}{G} (g_p + (g_p + g_{p-1})x + (g_p + g_{p-1} + g_{p-2})x^2 + \cdots + Gx^{p-1}). \end{aligned}$$

We use a classical result due to Eneström and Kakeya (see for example [8]): if $P(x) = \sum_{k=0}^n \alpha_k x^k$ is a (real) polynomial with all $\alpha_k > 0$, then the moduli of every root λ verifies

$$\min_{0 \leq k < n-1} \left\{ \frac{\alpha_k}{\alpha_{k+1}} \right\} \leq |\lambda| \leq \max_{0 \leq k < n-1} \left\{ \frac{\alpha_k}{\alpha_{k+1}} \right\}.$$

Applied to $R(x)$, we get

$$\min_{0 \leq k < p-1} \left\{ 1 - \frac{g_{p-k-1}}{\sum_{j=0}^{k+1} g_{p-j}} \right\} \leq |\lambda| \leq \max_{0 \leq k < p-1} \left\{ 1 - \frac{g_{p-k-1}}{\sum_{j=0}^{k+1} g_{p-j}} \right\}$$

hence the moduli of every root is at least $\frac{1}{G}$ and at most $\frac{G-1}{G}$, because Definition 1 specifies $g_1 \geq g_2 \geq \cdots \geq g_p > 0$. This is the first part of the fact.

For the second part, take a non null value λ of $Sp(DM)$ with an associated proper vector U_λ . Since $DM\mathbb{1}_p = 0_p$, the vectors U_λ and $\mathbb{1}_p$ are linearly independent. We have

$$\lambda U_\lambda = DMU_\lambda = MU_\lambda - \left(\frac{1}{G} (g_p, \dots, g_1) (MU_\lambda) \right) \mathbb{1}_p.$$

Let $\mu = \frac{1}{G} (g_p, \dots, g_1) (MU_\lambda)$. It follows that $MU_\lambda = \lambda U_\lambda + \mu \mathbb{1}_p$ and $M\mathbb{1}_p = \mathbb{1}_p$, thus

$$M((1-\lambda)U_\lambda - \mu\mathbb{1}_p) = (1-\lambda)(\lambda U_\lambda + \mu\mathbb{1}_p) - \mu\mathbb{1}_p = \lambda((1-\lambda)U_\lambda - \mu\mathbb{1}_p).$$

Thus, for $\lambda \neq 1$, the vector $(1-\lambda)U_\lambda - \mu\mathbb{1}_p$ is non null, from the linear independence of U_λ and $\mathbb{1}_p$, and is a proper vector of M for the eigenvalue λ .

For $\lambda = 1$, imagine that there exists a vector $U_1 = {}^t(u_1, u_2, \dots, u_p)$ such that $DMU_1 = U_1$. Since U_1 is in the image of DM , we have $\frac{1}{G} (g_p, \dots, g_1) U_1 = 0$.

Therefore $MU_1 = {}^t(u_2, \dots, u_p, 0)$. Thus, if we assume that $DMU_1 = U_1$, we necessarily have, looking at the last component,

$$U_1 = DMU_1 = MU_1 + u_p \mathbf{1}_p,$$

which gives, for $1 \leq i < p$, that $u_i = u_{i+1} + u_p$. Thus $U_1 = {}^t(pu_p, (p-1)u_p, \dots, u_p)$, and $(g_p, \dots, g_1)U_1 = u_p((g_p, \dots, g_1) {}^t(p, p-1, \dots, 1)) = 0$. The value $(g_p, \dots, g_1) {}^t(p, p-1, \dots, 1)$ is positive. Thus $u_p = 0$, and therefore $U_1 = \mathbf{0}_p$, which proves that 1 is not an eigenvalue of DM and finishes the proof. \square

Proof (Lemma 2). We recall Relation (2) of the perturbed weighted mean system,

$$\Delta v_{i+1} = m_i - \frac{\Delta h_{i+1}}{G}.$$

This proof strongly relies on the fact that the weighted mean is between the extremal values it is computed from,

$$\underline{m}_i \leq m_i \leq \overline{m}_i.$$

Furthermore, when $\underline{m}_i \neq \overline{m}_i$, this inequality is strict.

The difficult part is to show that the sequence $(\Delta v_k)_{k \geq i+d}$ is non-increasing. The proof is divided in two stages.

Fact 2 *There exists $d \leq (\overline{m}_i - \underline{m}_i + 1)p$, such that ΔV_{i+d} is a uniform vector, or $\Delta v_{i+d+1} < \underline{m}_{i+d}$.*

Proof (Fact 2). We prove that if the first case does not hold, then the second does. Assume that for each $d' \leq (\overline{m}_i - \underline{m}_i + 1)p$, $\Delta V_{i+d'}$ is not a uniform vector. In this case the mean $m_{i+d'}$ is always strictly between $\underline{m}_{i+d'}$ and $\overline{m}_{i+d'}$. We have $\overline{m}_{i+d'} > m_{i+d'} \geq \Delta v_{i+d'+1}$, thus the sequence $(\overline{m}_{i+d'})_{0 \leq d' \leq (\overline{m}_i - \underline{m}_i + 1)p}$ is non-increasing. Moreover for $d' \leq (\overline{m}_i - \underline{m}_i)p$, every p iterations the maximal value has strictly decreased: $\overline{m}_{i+d'} > \overline{m}_{i+d'+p}$, because the newly computed value (from the mean minus the discrete perturbation) is always strictly below and after p iterations we have only “newly competed values” so the maximal among them is strictly smaller. That is, we have $\overline{m}_{i+d'} \geq \overline{m}_{i+d'+p} + 1$ and, by induction, $\overline{m}_{i+p(\overline{m}_i - \underline{m}_i)} \leq \overline{m}_i - (\overline{m}_i - \underline{m}_i) = \underline{m}_i$. The non-uniformity also implies that $\underline{m}_{i+p(\overline{m}_i - \underline{m}_i)} < \underline{m}_i$. This last inequality is sufficient to conclude: in particular, there exists d'' , with $p(\overline{m}_i - \underline{m}_i) - p \leq d'' \leq p(\overline{m}_i - \underline{m}_i)$ such that $\Delta v_{i+d''} < \underline{m}_i \leq \Delta v_i$, which forces that there exists d with $0 \leq d < d''$ such that $\Delta v_{i+d+1} < \underline{m}_{i+d}$. \square

Fact 3 *The sequence $(\Delta v_k)_{k \geq i+d}$ is non-increasing.*

Proof (Fact 3). We first need the following preliminary inequality:

$$\text{for each integer } j > 0, m_{j+1} \leq m_j + \frac{g_1}{G}(\Delta v_{j+1} - \underline{m}_j).$$

Indeed,

$$\begin{aligned}
m_{j+1} &= \frac{1}{G} \sum_{k=1}^p g_k \Delta v_{j+2-k} \\
&= \frac{1}{G} \left(\sum_{k=2}^p (g_k - g_{k-1} + g_{k-1}) \Delta v_{j+2-k} + g_1 \Delta v_{j+1} \right) \\
&= \frac{1}{G} \left(\sum_{k=2}^p (g_k - g_{k-1}) \Delta v_{j+2-k} + \sum_{k=2}^p g_{k-1} \Delta v_{j+2-k} + g_1 \Delta v_{j+1} \right) \\
&= \frac{1}{G} \left(\sum_{k=2}^p (g_k - g_{k-1}) \Delta v_{j+2-k} + \sum_{k'=1}^{p-1} g_{k'} \Delta v_{j+1+k'} + g_1 \Delta v_{j+1} \right).
\end{aligned}$$

We have $g_p - g_1 \leq g_k - g_{k-1} \leq 0$ and $0 \leq \underline{m}_j \leq v_{j+2-k}$, thus

$$(g_p - g_1) \underline{m}_j \leq (g_k - g_{k-1}) \Delta v_{j+2-k}.$$

Moreover, $\sum_{k'=1}^{p-1} g_{k'} \Delta v_{j+1-k'} + g_p \Delta v_{j+1-p} = G m_j$. Thus

$$\begin{aligned}
m_{j+1} &\leq \frac{g_p - g_1}{G} \underline{m}_j + \left(m_j - \frac{g_p}{G} \Delta v_{j+1-p} \right) + \frac{g_1}{G} \Delta v_{j+1} \\
&= \frac{g_1}{G} (\Delta v_{j+1} - \underline{m}_j) + \frac{g_p}{G} (\underline{m}_j - \Delta v_{j+1-p}).
\end{aligned}$$

which gives the inequality announced since $\underline{m}_j \leq \Delta v_{j+1-p}$. In particular, we get $m_{j+1} \leq m_j$ when $\Delta v_{j+1} \leq \underline{m}_j$, and $m_{j+1} \leq m_j - \frac{g_1}{G}$ when $\Delta v_{j+1} < \underline{m}_j$.

Now we will prove by induction on k that, for each j such that $i + d \leq j \leq k$, we have $\Delta v_j \geq \Delta v_{j+1}$.

For initialization, this is true for $k = i + d$. Now assume that the property is true for an integer k , with $k \geq i + d$. This implies that the finite sequence $(m_j)_{i+d \leq j \leq k+1}$ is non-increasing and, for each j such that $i + d \leq j \leq k + 1$, we have $\Delta v_j = \underline{m}_j$.

If ΔV_{k+1} is a uniform vector, then $\Delta v_{k+2} \leq m_{k+1} = \Delta v_{k+1}$ and we are done. Otherwise let k' be the largest integer such that $k' \leq k$ and $\Delta v_{k'} > \Delta v_{k'+1}$ (notice that k' exists and is larger or equal to $k - p + 2$ because ΔV_{k+1} is not a uniform vector). Since $\frac{\Delta h_{k'+1}}{G} < 1 + \frac{g_1}{G}$ (Relation (5)), we have $m_{k'} - 1 - \frac{g_1}{G} < \Delta v_{k'+1}$. On the other hand, from the preliminary inequality we get $m_{k'+1} \leq m_{k'} - \frac{g_1}{G}$. Consequently

$$m_{k'+1} \leq m_{k'} - \frac{g_1}{G} < \Delta v_{k'+1} + 1 + \frac{g_1}{G} + -\frac{g_1}{G} = \Delta v_{k'+1} + 1,$$

which gives $m_{k'+1} < \Delta v_{k'+1} + 1$. Thus, since $(m_j)_{i+d \leq j \leq k+1}$ is non increasing, we have

$$\Delta v_{k+2} \leq m_{k'+1} < \Delta v_{k'+1} + 1 = \Delta v_{k+1} + 1$$

which gives $\Delta v_{k+2} \leq \Delta v_{k+1}$, since Δv_{k+2} and Δv_{k+1} are integer values. This finishes the induction. \square

To finish the proof of Lemma 2, we have

$$\Delta v_{k+1} > m_k - 1 - \frac{g_1}{G} \geq m_k - 2 \geq \underline{m}_k - 2 = \Delta v_k - 2$$

which gives $\Delta v_{k+1} \geq \Delta v_k - 1$, by integrity. \square

Proof (Lemma 3). From the hypothesis, we have

$$\Delta h_{n_2+k+1} = m(\Delta^2 V_{n_2+k}) + \Delta^2 v_{n_2+k} G \leq m(1, 1, \dots, 1) < g_1.$$

We have $\Delta^2 V_{n_2+k+r-1} \leq (1, 1, \dots, 1) - E_{p+1-r}$, which gives

$$\Delta h_{n_2+k+r} = m(\Delta^2 V_{n_2+k+r-1}) + \Delta^2 v_{n_2+k+r-1} G \leq m(1, 1, \dots, 1) - m(E_{p+1-r}) + G.$$

Notice that $m(E_{p+1-r}) = g_r + g_{r+1} + \dots + g_p$. As a consequence

$$\Delta h_{n_2+k+r} \leq m(1, 1, \dots, 1) - (g_r + g_{r+1} + \dots + g_p) + G = m(1, 1, \dots, 1) + g_1 + g_2 + \dots + g_{r-1}.$$

Thus

$$\Delta h_{n_2+k+r} + g_r \leq m(1, 1, \dots, 1) + g_1 + g_2 + \dots + g_r < g_1 + G$$

which, by Theorem 3 (item 3), ensures that $a_j = 0$ for $j \geq n_2 + k + 1$. \square

Proof (Theorem 4). The proof proceeds by induction on N . For $N = 0$, the result is obvious. Now assume that the result is true for $N \geq 0$. By induction hypothesis, the sequence $(\Delta^2 v_i)_{i \geq n_2}$ is element of $1^* (0 + \epsilon) 1^* 0^\omega$. If $(\Delta^2 a_i)_{i \geq n_2} = 0^\omega$, then $(\Delta^2 v'_i)_{i \geq n_2} = (\Delta^2 v_i)_{i \geq n_2}$, and the result is trivial. Otherwise $(\Delta^2 a_i)_{i \geq n_2}$ contains at least a value equal to 1. We have three alternatives. The key point is that a pattern $(-1)1$ in $\Delta^2 a$ must occur on a pattern 10 in $\Delta^2 v$.

– If $(\Delta^2 v_i)_{i \geq n_2} = 1^k 0^\omega$ for an integer k , then the last 1 of $(\Delta^2 a_i)_{i \geq n_2}$ must be in a position greater than k , since, otherwise, a value 2 appears in $(\Delta^2 v'_i)_{i \geq n_2}$ by addition. Moreover the last -1 of $(\Delta^2 a_i)_{i \geq n_2}$ must be in a position at most k since, otherwise, a value -1 appears in $(\Delta^2 v'_i)_{i \geq n_2}$ by addition (Equation (6)). Since the last -1 of $(\Delta^2 a_i)_{i \geq n_2}$ just precedes the last 1 of $(\Delta^2 a_i)_{i \geq n_2}$, the last 1 of $(\Delta^2 a_i)_{i \geq n_2}$ is in position $k + 1$ and the last -1 of $(\Delta^2 a_i)_{i \geq n_2}$ is in position k .

If $(\Delta^2 a_i)_{i \geq n_2} = 0^{k-1} (-1)1 0^\omega$, then, by addition, $(\Delta^2 v'_i)_{i \geq n_2} = 1^{k-1} 01 0^\omega$, and we are done. Otherwise, we study the position of the penultimate -1 of $(\Delta^2 a_i)_{i \geq n_2}$. Since this -1 is just preceded by a 1, the only non contradictory possibility is that -1 is in first position *i.e.*, $\Delta^2 a_{n_2} = -1$. Thus $\Delta^2 a_{n_2-1} = 1$, which forces that $\Delta^2 v_{n_2-1} = 0$, since $\Delta^2 v'_{n_2-1} = \Delta^2 v_{n_2-1} + \Delta^2 a_{n_2-1} \leq 1$. As a consequence, from Lemma 3, $\Delta^2 a_j = 0$ for $j > n_2$, which is a contradiction and finishes the proof of this case.

– If $(\Delta^2 v_i)_{i \geq n_2} = 1^k 0 1^{k'} 0^\omega$, for a pair k, k' of positive integers, then from Lemma 3 the avalanche stops before column $n_2 + k$ *i.e.*, $a_j = 0$ for $j > n_2 + k$. Thus, proceeding as is the previous item, we get that the last 1 of $(\Delta^2 a_i)_{i \geq n_2}$ must be in position $k + 1$, from which we deduce that $(\Delta^2 a_i)_{i \geq n_2} = 0^{k-1} (-1)1 0^\omega$. Thus, by addition, $(\Delta^2 v'_i)_{i \geq n_2} = 1^{k-1} 0 1^{k'+1} 0^\omega$, and we are done.

- The case when $(\Delta^2 v_i)_{i \geq n_2} = 0 1^{k'} 0^\omega$, with $k' > 0$ can be treated as the previous one. We have $(\Delta^2 a_i)_{i \geq n_2} = 1 0^\omega$. Thus, by addition, $(\Delta^2 v'_i)_{i \geq n_2} = 1^{k'+1} 0^\omega$, and we are done.

□

Proof (Lemma 4). Let $q = q_1 q_2 \dots q_{p-1}$ be a state of $\mathcal{L}(\mathcal{A}'_{\mathcal{R}})$. We prove the following equivalent statement: if $q_j = 1$, then, for $1 \leq k \leq p-1-r$ and $j+k \leq p-1$, we have $q_{j+k} = 0$. For this purpose, we demonstrate that after i automaton steps, the result holds for each $j \geq p-1-i$.

For $i = 0$, the result is obvious. For the induction, assume that the result holds for a fixed $i \geq 0$. Let $q_1 q_2 \dots q_{p-1}$ be a state obtained after $i+1$ steps, whose predecessor is the state $q_0 q_1 \dots q_{p-2}$. Since $q_0 q_1 \dots q_{p-2}$ is obtained after i steps, we can apply the induction hypothesis on $q_0 q_1 \dots q_{p-2}$: if $q_j = 1$ and $j \geq p-2-i$, then, for $1 \leq k \leq p-r$ and $j+k \leq p-2$ we have $q_{j+k} = 0$. If, moreover, $q_{p-1} = 0$, then we obviously have the result on $q_1 q_2 \dots q_{p-1}$ for $i+1$. Otherwise $q_{p-1} = 1$, which forces that $m(q_0 q_1 \dots q_{p-2}) < g_1$, since there is an automaton transition from $q_0 q_1 \dots q_{p-2}$ to $q_1 q_2 \dots q_{p-1}$. Therefore, for any j , with $0 \leq j \leq p-2$, $m(E_{j+1}) \leq m(q_0 q_1 \dots q_{p-2}) < g_1$. Furthermore for any j with $r \leq j \leq p-2$ we have $q_j = 0$, since $m(E_{j+1}) \geq m(E_{r+1}) > g_1$. As a consequence, there is nothing to prove for $j \geq r$. For $p-i-2 \leq j < r$, the condition is satisfied by induction hypothesis. This gives the result by induction.

□

Proof (Lemma 5). First notice that our conditions and Lemma 4 imply that $\Delta^2 v_{j-k} = 0$ for $1 \leq r \leq p-r-1$, and $\Delta^2 V_j$ contains at most one non null component. Let k_0 be the lowest positive integer k such that $\Delta^2 v_{j+k} = 1$ (with the classical convention $k_0 = +\infty$ if $\Delta^2 v_{j+k} = 0$ for any $k > 0$). We have $k_0 \geq p-r+1$ from Lemma 4.

The remarks above allow us to apply the criterion of Theorem 3 (item 3). Indeed, for $j+2$ we have

$$\Delta h_{j+2} + g_1 = m(\Delta^2 V_{j+1}) + g_1 \leq m(E_{p-1}) + m(E_{r-1}) + g_1.$$

Furthermore $m(E_{p-1}) + g_1 = G$ and $m(E_{r-1}) < m(E_r) \leq g_1$, thus

$$\Delta h_{j+2} + g_1 < G + g_1.$$

More generally, for $2 \leq k \leq k_0$ we have

$$\Delta h_{j+k} + g_{k-1} = m(\Delta^2 V_{j+k-1}) + g_{k-1} \leq m(\Delta^2 V_{j+1}) + g_1 = \Delta h_{j+2} + g_1 < G + g_1.$$

Now, for $k = k_0 + 1$ we have

$$\Delta h_{j+k_0+1} + g_{k_0} = m(\Delta^2 V_{j+k_0}) + G + g_{k_0}.$$

Remark that $m(\Delta^2 V_{j+k_0}) = m(E_{p-1-k_0})$, thus

$$m(\Delta^2 V_{j+k_0}) + g_{k_0} = g_p + g_{p-1} + \dots + g_{k_0} \leq m(E_{p-k_0}) \leq m(E_{r-1}) < g_1,$$

and we can deduce that

$$\Delta h_{j+k_0+1} + g_{k_0} < G + g_1.$$

Finally, for $k_0 + 2 \leq k \leq p + 1$ we have $\Delta^2 v_{j+k-1} = 0$ from Lemma 4. Moreover, $\Delta^2 V_{j+k-1} = E_{p-1-k+k_0} + E_{p-1-k}$ with $m(E_{p-1-k+k_0}) < G$ and $m(E_{p-1-k}) + g_{k-1} = g_p + g_{p-1} + \cdots + g_{k-1} \leq m(E_{p-k}) \leq m(E_{r-1}) < g_1$. Hence we can again deduce that

$$\Delta h_{j+k} + g_{k-1} < G + g_1.$$

As a conclusion, from Theorem 3 (item 3) we can claim that $(a_i)_{i \geq j+2} = 0^\omega$. \square

Proof (Theorem 5). By induction on N . For $N = 0$, the result is obvious. Now assume that the result holds for a fixed $N \geq 0$. If $(a_i)_{i > n_2+p} = 0^\omega$, then the result is straightforward. Otherwise we have two cases.

- If $(\Delta^2 v_i)_{i > n_2+p}$ is a suffix of an element of $(0^{p-r-1} 1)^* 0^\omega$, then the position of the last -1 of $(a_i)_{i > n_2+p}$ (just before the last 1 of $(a_i)_{i > n_2+p}$) and the position of the last 1 of $(\Delta^2 v_i)_{i > n_2+p}$ are the same. Otherwise, by addition (Equation (6)), a value 2 or -1 appears in $(\Delta^2 v'_i)_{i > n_2+p}$, or two consecutive 1 are separated by $p - r - 2$ values 0 in $(\Delta^2 v'_i)_{i > n_2+p}$. The first alternative is in contradiction with Proposition 1, and the second one with Lemma 4. For the same reasons, the position of the penultimate -1 of $(a_i)_{i \geq 0}$ is before $n_2 + p$. This gives the result for $N + 1$, by addition.
- If $(\Delta^2 v_i)_{i > n_2+p}$ is a suffix of an element of $(0^{p-r-1} 1)^* 0 (0^{p-r-1} 1)^* 0^\omega$, then the position of the last -1 of $(a_i)_{i > n_2+p}$ and the position of the 1 just before the $p - r$ values 0 in $(\Delta^2 v_i)_{i > n_2+p}$ are the same. Otherwise, we have the same contradictions as in the previous case. Moreover, we know from Lemma 5 that the last -1 in $(a_i)_{i > n_2+p}$ cannot be after the position of the 1 just before the $p - r$ values 0 in $(\Delta^2 v_i)_{i > n_2+p}$. This ensures that the last -1 of $(a_i)_{i > n_2+p}$ has not the same position as the last 1 of $(\Delta^2 v_i)_{i > n_2+p}$. Continuing in a similar way as in the previous case, we get the result for $N + 1$.

This proves the lemma by induction. \square