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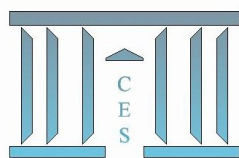
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On integer-valued means and the symmetric maximum

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Abstract

Integer-valued means, satisfying the decomposability condition of Kolmogoroff/Nagumo, are necessarily extremal, i.e., the mean value depends only on the minimal and maximal inputs. To overcome this severe limitation, we propose an infinite family of (weak) integer means based on the symmetric maximum and computation rules. For such means, their value depends not only on extremal inputs, but also on 2nd, 3rd, etc., extremal values as needed. In particular, we show that this family can be characterized by a weak version of decomposability.

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Keywords: integer means; nonassociative algebra; symmetric maximum; decomposability

1 Introduction

In 1930, Kolmogoroff [5] and Nagumo [7] independently characterized the class of quasi-arithmetic means (including the well-known arithmetic, geometric and harmonic means) through a property which was later called *decomposability* [4], *barycentric associativity* [6], or simply *associativity* [1]. As associativity in the classical algebraic sense will play a central rôle in this paper, we prefer to adopt the term “decomposability”.

Decomposability is a very natural property that amounts to saying that the average of a set of numbers should not be changed if a subset of it is replaced by its average. Denoting such an averaging function by μ , the latter notion can be expressed by the condition

$$\mu(a_1, \dots, a_n) = \mu(\underbrace{b, \dots, b}_{k \text{ times}}, a_{k+1}, \dots, a_n) \quad (1)$$

where $b = \mu(a_1, \dots, a_k)$.

More recently, Bennett et al. [1] investigated the effect of this property on integer-valued averaging functions, that is, averaging functions defined and valued in \mathbb{Z} . This setting is both natural and relevant in practice since in many situations input values are integer, for instance, number of children, of items, etc., and in such cases mean values that are not integer bear no meaning. Surprisingly, by imposing decomposability in addition to the usual requirements of averaging functions (nondecreasingness, anonymity and internality), it turns out that such integer-valued averaging functions (called *integer means* or *\mathbb{Z} -means* in [1]) are *extremal*, that is, the output depends only on the minimum and maximum input values. In symbols,

$$\mu(a_1, \dots, a_n) = \mu(a_1, a_n)$$

with $a_1 \leq a_2 \leq \dots \leq a_n$. Put otherwise, integer means ignore all input values but the minimum and the maximum ones.

The aim of this paper is to attempt to finding a class of integer-valued averaging functions that have a less drastic behavior and whose output retains more information of the input, by allowing to considering not only minimal and maximal inputs, but also, when necessary, second, third, etc., min/max values. This can be achieved through the use of the *symmetric maximum*, a symmetrization of the classical maximum operator proposed by Grabisch [3], and denoted by \oplus . Roughly speaking, \oplus is a binary operator on \mathbb{Z} yielding the extremal value of its inputs (i.e., the greatest in absolute value), except if these are opposite values, in which case the output is 0, for instance

$$a \oplus (-a) = 0, \quad 2 \oplus 3 = 3, \quad (-2) \oplus 3 = 3, \quad 2 \oplus (-3) = -3.$$

As it is easy to see, such an operator is not associative in the algebraic sense. Indeed,

$$(3 \oplus (-3)) \oplus 2 = 0 \oplus 2 = 2, \quad 3 \oplus ((-3) \oplus 2) = 3 \oplus (-3) = 0. \quad (2)$$

Therefore, defining an averaging function by \oplus , i.e., writing $\oplus(a_1, \dots, a_n)$ for any sequence a_1, \dots, a_n of inputs is meaningless without further refinement of the definition. In fact, nonassociativity arises if and only if there exist extremal opposite terms in the input sequence (like 3, -3), otherwise $\oplus(a_1, \dots, a_n)$ is well defined. In order to overcome this drawback, Grabisch [3] proposed systematic ways to aggregate all terms of a sequence using a single binary operator as in (2). These systematic ways of putting parentheses were called *rules of computation*, and were fully investigated by the authors in [2]. A simple example of such a rule is the following: aggregate first all pairs of extremal opposite numbers in decreasing order, till no such pair exist. On the sequence of number 4, -4, 3, -3, 2, 1, this would yield:

$$\oplus(4, -4, 3, -3, 2, 1) = (4 \oplus (-4)) \oplus (3, -3, 2, 1) = 0 \oplus (((3 \oplus (-3)) \oplus (2, 1)) = 0 \oplus (2, 1) = 2.$$

\oplus viewed as an integer-valued averaging function is clearly not extremal, but as said above, it satisfies extremality in a wider sense: in the above example we have

$$\oplus(4, -4, 3, -3, 2, 1) = \oplus(2, 1)$$

and (2,1) is the 3d pair of min/max inputs, the two first pairs having been eliminated because they are pairs of opposite values.

As shown in [2], infinitely many rules of computation exist, yielding an infinite (and even uncountable) family of integer-valued averaging functions. Interestingly, none of them satisfy decomposability (Corollary 1), but they satisfy a weaker version of it, hence we call them *weak integer means*. More precisely, the decomposability property (1) is valid only for certain subsequences of a_1, \dots, a_n (Lemma 2).

We provide a complete characterization of this family of weak integer means, by using weak decomposability (Theorem 3), or a related property called the replacement property (Lemma 5). The various characterizations are summarized in Theorem 4.

2 Preliminaries

2.1 \mathbb{Z} -means and the symmetric maximum

We consider a countable set C which is totally ordered by \leq , and with minimal element 0. Let $C^- = \{-c \mid c \in C\}$ be its dually isomorphic copy, and consider $\tilde{C} = C \cup C^-$, setting $-0 = 0$.¹ For convenience, we introduce the notation

$$\tilde{C}^* = \left(\bigcup_{n \in \mathbb{N}} (\tilde{C})^n \right) \cup \{\varepsilon\},$$

where ε is the empty string. Put differently, \tilde{C}^* is the set of all words built over the alphabet \tilde{C} using concatenation.

Throughout the paper we deal with functions $\mu : \tilde{C}^* \rightarrow \tilde{C}$, and always adopt the following convention: $\mu(\varepsilon) = 0$ and $\mu(a) = a$ for all $a \in \tilde{C}$.

We recall the main result from [1]. A function

$$\mu : \tilde{C}^* \rightarrow \tilde{C}$$

is a \mathbb{Z} -mean or *integer-valued mean* if it satisfies the following properties:

- (i) *Anonymity*: for every $n \in \mathbb{N}$, every permutation π on $[n]$,

$$\mu(a_1, \dots, a_n) = \mu(a_{\pi(1)}, \dots, a_{\pi(n)});$$

- (ii) *Internality*: for every $n \in \mathbb{N}$, if $a_1 \leq \dots \leq a_n$,

$$a_1 \leq \mu(a_1, \dots, a_n) \leq a_n;$$

- (iii) *Monotonicity*: for every $n \in \mathbb{N}$, $a_i \leq a'_i$ for $i = 1, \dots, n$ implies

$$\mu(a_1, \dots, a_n) \leq \mu(a'_1, \dots, a'_n);$$

¹A typical example is $\tilde{C} = \mathbb{Z}$, and in fact any such \tilde{C} is isomorphic to \mathbb{Z} or to a subset of it. This is why we may call elements of \tilde{C} “integers” and use the corresponding terminology, like positive and negative numbers, absolute value, etc.

(iv) *Decomposability*²: for every $k, n \in \mathbb{N}$ with $k \leq n$ and $\mu(a_1, \dots, a_k) = b$, we have

$$\mu(a_1, \dots, a_k, a_{k+1}, \dots, a_n) = \mu(\underbrace{b, \dots, b}_k, a_{k+1}, \dots, a_n).$$

A function $\mu : \tilde{C}^* \rightarrow \tilde{C}$ is *extremal* if for any $n \in \mathbb{N}$ and any $a_1 \leq \dots \leq a_n$, we have

$$\mu(a_1, \dots, a_n) = \mu(a_1, a_n)$$

Theorem 1. All \mathbb{Z} -means are extremal.

Hence, \mathbb{Z} -means, unlike means defined on real numbers, do not consider all values to be aggregated but only the extreme ones. Our aim is to overcome this rather negative result by weakening the definition of decomposability, which is mainly responsible for extremality, and to propose a whole family of aggregation functions fulfilling decomposability in a weak sense. This will be achieved through the *symmetric maximum* [3], which is a binary operation $\oslash : \tilde{C}^2 \rightarrow \tilde{C}$ defined by

$$a \oslash b = \begin{cases} -(|a| \vee |b|), & \text{if } b \neq -a \text{ and } |a| \vee |b| = -a \text{ or } -b \\ 0, & \text{if } b = -a \\ |a| \vee |b|, & \text{otherwise.} \end{cases} \quad (3)$$

In other words, except if $a = -b$, $a \oslash b$ equals the element among a, b that is *extremal*, i.e., that has the largest absolute value (magnitude). It is not difficult to see that \oslash satisfies the following properties:

- (i) (C1) \oslash coincides with the maximum on C^2 ;
- (ii) (C2) $a \oslash (-a) = 0$ for every $a \in \tilde{C}$;
- (iii) (C3) $-(a \oslash b) = (-a) \oslash (-b)$ for every $a, b \in \tilde{C}$.

By (C1), (C3), \oslash coincides with the minimum on $(C^-)^2$. Also, it is monotone, internal, commutative (or symmetric), and 0 is its neutral element (by (C1) and (C3)). Hence, \oslash almost behaves like a commutative group operation on \tilde{C} , however it lacks associativity: it is not true in general that

$$a \oslash (b \oslash c) = (a \oslash b) \oslash c, \quad (4)$$

as shown by the following example:

$$(-3 \oslash 3) \oslash 1 = 0 \oslash 1 = 1; \quad -3 \oslash (3 \oslash 1) = -3 \oslash 3 = 0.$$

The following result gives an algebraic characterization of \oslash and shows that if one requires that (C1), (C2) and (C3) hold, then (3) is the best possible definition for \oslash .

Theorem 2. [3, Prop. 5] No binary operation satisfying (C1), (C2), (C3) fulfills (4) on a larger domain than \oslash .

²This property is called associativity in [1], and barycentric associativity in [6]. However, the usual meaning of associativity in algebra is different, and the latter is a central property in this paper. To avoid any confusion, we therefore stick to the term “decomposability”.

It is easy to see that the domain where \mathbb{V} is associative is given by the set of triples (a, b, c) such that

$$\min(a, b, c) \neq -\max(a, b, c). \quad (5)$$

Theorem 2 states that under (C1), (C2) and (C3), any binary operation μ on \tilde{C} is nonassociative as soon as (5) is not satisfied (i.e., for triples with opposite extremal elements). Moreover, if $\mu \neq \mathbb{V}$, then there exists a triple satisfying (5) for which μ is nonassociative. For instance, $\overset{0}{\mathbb{V}}$ defined by

$$\overset{0}{\mathbb{V}}(a_1, \dots, a_n) = \begin{cases} a_1, & \text{if } a_n \leq 0 \\ 0, & \text{if } a_1 < 0 < a_n \\ a_n, & \text{if } a_1 \geq 0 \end{cases} \quad (6)$$

with $a_1 \leq \dots \leq a_n$, satisfies (C1), (C2) and (C3), but

$$(-2 \overset{0}{\mathbb{V}} 2) \overset{0}{\mathbb{V}} 3 = 0 \overset{0}{\mathbb{V}} 3 = 3; \quad -2 \overset{0}{\mathbb{V}} (2 \overset{0}{\mathbb{V}} 3) = -2 \overset{0}{\mathbb{V}} 3 = 0.$$

2.2 How to make \mathbb{V} associative: Rules of computation

The lack of associativity implies that it is not possible to extend the definition of \mathbb{V} to \tilde{C}^* , since the result depends on the particular way that we apply the binary \mathbb{V} to the terms of a sequence with more than two terms. Therefore, no function like a \mathbb{Z} -mean can be defined based on this operation. However, a way to overcome this difficulty, suggested in [3] and fully developed in [2], is to fix beforehand a systematic way of putting parentheses on any sequence of \tilde{C}^* , procedure that was called a *rule of computation*.

Elements of \tilde{C}^* are called sequences (of “integers”), denoted by $\sigma = (a_i)_{i \in I}$ for some finite index set I , and $a_i \in \tilde{C}$ for all $i \in I$ (note that 0 is a possible value, contrarily to [2]). Since we consider only anonymous functions on \tilde{C}^* , the order is unimportant and so we consider only sequences of integers in decreasing order of their absolute value, positive integers first (e.g., 5, 5, -5, -3, 2, -2, 1, 0). The empty sequence is ε . The set of all such integer sequences, including the empty sequence, is denoted by \mathfrak{S} . Since sequences are ordered as described above, a convenient notation is the following. For an arbitrary sequence

$$\sigma = (\underbrace{n_1, \dots, n_1}_{p_1 \text{ times}}, \underbrace{-n_1, \dots, -n_1}_{m_1 \text{ times}}, \dots, \underbrace{n_q, \dots, n_q}_{p_q \text{ times}}, \underbrace{-n_q, \dots, -n_q}_{m_q \text{ times}})$$

we introduce $\theta(\sigma) = (n_1, \dots, n_q)$, the sequence of absolute values (magnitudes) of integers in σ , and $\psi(\sigma) = ((p_1, m_1), \dots, (p_q, m_q))$ all numbers of occurrence of these integers.

Example 1. Consider the sequence $\sigma = (3, 3, -3, 2, -2, -2, 1, 1, 1, 1)$. Then

$$\theta(\sigma) = (3, 2, 1); \quad \psi(\sigma) = ((2, 1), (1, 2)(4, 0)).$$

A sequence $(a_i)_{i \in I}$ fulfills *associativity* (or, simply, is *associative*) if $\min_{i \in I} a_i \neq -\max_{i \in I} a_i$ (equivalently, if either $p_1 = 0$ or $m_1 = 0$). The set of sequences which do not fulfill associativity is denoted by \mathfrak{S}_0 . Note that by this definition the sequence $(1, -1)$, although containing only 2 elements, is not associative.

The above notation permits to introduce rules of computation in a proper way. We define 5 elementary rules $\rho_i : \mathfrak{S} \rightarrow \mathfrak{S}$ which operate on σ as follows:

- (i) Elementary rule ρ_1 : if $p_1 > 1$ and $m_1 > 0$, then p_1 is changed to $p_1 = 1$;
- (ii) Elementary rule ρ_2 : idem with p_1, m_1 exchanged;
- (iii) Elementary rule ρ_3 : if $p_1 > 0$, $m_1 > 0$, the pair (p_1, m_1) is changed into $(p_1 - c, m_1 - c)$, where $c = p_1 \wedge m_1$;
- (iv) Elementary rule ρ_4 : if $p_1 > 0$, $m_1 > 0$, and if $p_2 > 0$, then p_2 is changed into $p_2 = 0$;
- (v) Elementary rule ρ_5 : idem with m_2 replacing p_2 .

Hence, elementary rules delete terms in sequences which do not fulfill associativity, and do nothing otherwise.

A (*well-formed*) *computation rule* R is a word built with the alphabet $\{\rho_1, \dots, \rho_5\}$, i.e., $R \in \mathcal{L}(\rho_1, \dots, \rho_5)$, such that $R(\sigma) \in \mathfrak{S} \setminus \mathfrak{S}_0$ for all $\sigma \in \mathfrak{S}$. The set of (well-formed) computation rules is denoted by \mathfrak{R} . Examples of rules are (words are read from left to right)

- (i) $\langle \cdot \rangle_0 = \rho_3^*$;
- (ii) $\langle \cdot \rangle_- = (\rho_1 \rho_2 \rho_3)^*$;
- (iii) $\langle \cdot \rangle_-^+ = (\rho_4 \rho_5)^* \rho_1 \rho_2 \rho_3$.

We use the notation \oplus_R for $\oplus \circ R$. It is shown in [2] that any computation rule has the form $R = T_1 T_2 \dots$, where each T_i has the form $\omega \rho_1^\alpha \rho_2^\beta \rho_3$, with $\omega \in \mathcal{L}(\rho_4, \rho_5)$ and $\alpha, \beta \in \{0, 1\}$ (factorization scheme).³ Moreover, \mathfrak{R} is uncountable (the set of its maximal elements is isomorphic to $2^{\mathbb{N}}$). We denote by \mathfrak{R}_{123} the set of rules where $\omega = \varepsilon$ in each term T_i (i.e., the rule uses only ρ_1, ρ_2, ρ_3). Also, we introduce $\text{Ker}(R) = \{\sigma \in \mathfrak{S} \mid R(\sigma) = \varepsilon\}$, the *kernel* of R .

Two rules R, R' are equivalent (denoted by $R \sim R'$) if $\oplus_R = \oplus_{R'}$ on \mathfrak{S} . For a given rule R and a sequence $\sigma = (a_i)_{i \in I}$, we denote by J_σ^R the set of elements deleted by R in σ . We use the following notation: for a sequence $(a_i)_{i \in I}$ and any set $K \subseteq I$, we write a_K for the subsequence $(a_i)_{i \in K}$. Hence, we may write

$$(a_i)_{i \in I} = (a_K, a_{I \setminus K}).$$

For convenience, we may also write σ_K for a restriction of the sequence σ to K .

2.3 The basic mechanism and properties of \oplus_R

It is important to analyse the effect of a rule $R = T_1 T_2 \dots$ (not necessarily in \mathfrak{R}) on a sequence $\sigma \in \mathfrak{S}_0$, with representation

$$\psi(\sigma) = ((p_1, m_1), \dots, (p_r, m_r)), \quad \theta(\sigma) = (n_1, \dots, n_r).$$

Let us first consider a single term $T_1 = \omega \rho$, where ρ is a shorthand for $\rho_1^\alpha \rho_2^\beta \rho_3$. It is straightforward to verify that $T_1(\sigma)$ is another sequence with representation

$$\psi(T_1(\sigma)) = ((p'_1, m'_1), (p'_q, m'_q), (p_{q+1}, m_{q+1}), \dots, (p_r, m_r)), \quad \theta(T_1(\sigma)) = (n_1, n_q, \dots, n_r)$$

³Here, we adopt the notation $\rho^0 = \varepsilon$ and $\rho^1 = \rho$.

with the following requirements:

$$\begin{aligned} 0 \leq p'_1 < p_1, \quad 0 \leq m'_1 < m_1, \text{ and at most one of them is positive} \\ 2 \leq q \leq r \\ p'_q = 0 \text{ or } p_q, \quad m'_q = 0 \text{ or } m_q \end{aligned}$$

This situation is pictured in Figure 1. The first hatched block is the part deleted by

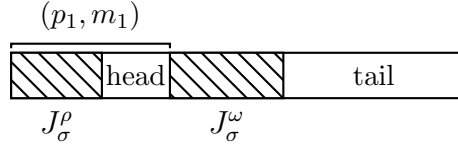


Figure 1: The sequences σ and $T_1(\sigma)$ (white part), with values from left to right in decreasing order of their magnitude. The hatched part is $J^{T_1}_\sigma$

ρ , while the second one is deleted by ω . The “head” is the undeleted part of (p_1, m_1) between the two hatched blocks and it may be the case that it does not exist. It contains only n_1 or only $-n_1$, with possible repetitions. The “tail” is the part left undeleted by ω and may not exist. The left hatched block always exist, while the second may not, in which case we say that the head is empty.

Applying successively the different terms T_2, T_3, \dots of R it is not difficult to see that the final result $R(\sigma)$ has almost the same form. Indeed, if a head exists in $T_1(\sigma)$, then $T_1(\sigma)$ is associative, and the rule stops. If not, the next term applies (if nonassociative), yielding the same form as above. Hence the representation for $R(\sigma)$ is

$$\psi(R(\sigma)) = ((p'_k, m'_k), (p'_q, m'_q), (p_{q+1}, m_{q+1}), \dots, (p_r, m_r)), \quad \theta(R(\sigma)) = (n_k, n_q, \dots, n_r)$$

with the following requirements:

$$\begin{aligned} k \geq 1 \\ 0 \leq p'_k \leq p_k, \quad 0 \leq m'_k \leq m_k, \text{ and at most one of them is positive} \\ k+1 \leq q \leq r \\ p'_q = 0 \text{ or } p_q, \quad m'_q = 0 \text{ or } m_q \end{aligned}$$

This situation is pictured in Figure 2. We define the head and tail of $R(\sigma)$ similarly:

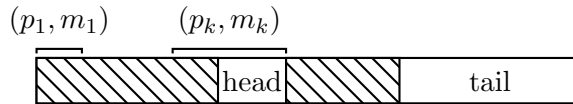


Figure 2: The sequences σ and $R(\sigma)$ (white part), with values from left to right in decreasing order of their magnitude. The hatched part is J^R_σ

the head is the part left undeleted by R between the two deleted parts, while the tail is

the part left undeleted which has no deleted part on its right side. If the head exists, it contains only value n_k or only $-n_k$, possibly with repetition. Note that when they both exist, the elements in the tail have strictly smaller magnitude than any element of the head.

We denote by H, T the subsets of I corresponding to the head and tail of a sequence $\sigma = (a_i)_{i \in I}$, and therefore may use the notation σ_H, σ_T or a_H, a_T for the head and tail.

Based on this result, we make some useful observations.

Lemma 1. Let $\sigma = (a_i)_{i \in I} \in \mathfrak{S}_0$. The following holds.

- (i) $\mathbb{O}_R(\sigma) = \mathbb{O}_R(a_{I \setminus J_\sigma^R}) = \mathbb{O}(a_{I \setminus J_\sigma^R}) = \left(\min_{i \in I \setminus J_\sigma^R} a_i \right) \mathbb{O} \left(\max_{i \in I \setminus J_\sigma^R} a_i \right)$ for any $R \in \mathfrak{R}$;
- (ii) \mathbb{O}_R is odd for any $R \in \mathfrak{R}$, i.e., $\mathbb{O}_R(-\sigma) = -\mathbb{O}_R(\sigma)$;
- (iii) Let $R = T_1 T_2 \dots$. If $R(\sigma)$ has a (nonempty) head, then $\mathbb{O}_R(\sigma_{I \setminus T}, \sigma_T) = \mathbb{O}_R(\sigma_{I \setminus T}, \sigma')$, where σ' is any sequence (possibly empty); If $R(\sigma)$ has no head, then $\mathbb{O}_R(\sigma_{I \setminus T}, \sigma_T) = \mathbb{O}_R(\sigma_{I \setminus T}, \sigma')$, for any $\sigma' \in \mathfrak{S} \setminus \mathfrak{S}_0$ which has the same extremal value as σ_T ;
- (iv) $a_{J_\sigma^R} \in \text{Ker}(R)$, hence $\mathbb{O}_R(a_{J_\sigma^R}) = 0$, for any R of the form $T_1 T_2 \dots$ (finite or infinite);
- (v) For any $R \in \mathfrak{R}$, letting $\mathbb{O}_R(\sigma) = b$, then $|b| \leq \max_{i \in J_\sigma^R} |a_i|$, and $|b| \geq |a_i|$ for all $i \in I \setminus J_\sigma^R$ (and there is equality for some i , if $I \setminus J_\sigma^R \neq \emptyset$);

Proof. (i) The first equality comes from $\mathbb{O}_R(\sigma) = \mathbb{O}(R(\sigma)) = \mathbb{O}(a_{I \setminus J_\sigma^R})$. By definition of a computation rule, $a_{I \setminus J_\sigma^R}$ fulfills associativity, which explains the second equality, while the third one follows from associativity and the definition of \mathbb{O} .

(ii) From (i) and (C3).

(iii) Suppose $R(\sigma)$ has a head. Then $\mathbb{O}_R(\sigma) = b$ is the value in the head, which is strictly greater in magnitude than any value in the tail. Therefore the tail has no effect on the result. Otherwise, b is the extremal value of the tail. Hence the tail can be modified provided that it remains associative and with same extremal value.

(iv) Step 1: We consider first that $R = T_1 = \omega \rho_1^\alpha \rho_2^\beta \rho_3$ and we denote $\rho = \rho_1^\alpha \rho_2^\beta \rho_3$ for short. Consider a sequence $\sigma = (a_i)_{i \in I}$ and let $\psi(\sigma) = ((p_1, m_1), \dots, (p_k, m_k))$, $\psi(a_{J_\sigma^R}) = ((p'_1, m'_1), \dots, (p'_{k'}, m'_{k'}))$. Observe that $p'_1 > 0$, $m'_1 > 0$ since ρ_3 is present in T_1 . Therefore, if $\varepsilon \neq \omega = \omega_1 \omega_2 \dots$, ω_1 applies, then the remaining ω_2 , etc., also apply because no ω_i modifies p'_1, m'_1 . It then follows that $J_\sigma^\omega = J_{a_{J_\sigma^R}}^\omega$. Now, observe that ρ modifies only (p_1, m_1) , hence $J_{\omega(\sigma)}^\rho = J_{\omega(a_{J_\sigma^R})}^\rho$, and thus $R(\sigma_{J_\sigma^R}) = \varepsilon$.

Step 2: We consider now the case where R is not limited to a single term T_1 and we show the result by induction on the size of the sequence.

Induction hypothesis: for any $\sigma = (a_i)_{i \in I} \in \mathfrak{S}_0$, $|\sigma| \leq t$ with $t \geq 3$, for any $R = T_1 T_2 \dots$, $R(a_{J_\sigma^R}) = \varepsilon$.

Step 2.1: We show the base case with $t = 3$. Any sequence in \mathfrak{S}_0 of three terms is of the form $(a, -a, b)$ with $|a| \geq |b|$. Then only T_1 is applied in R , and it is easy to check that the result holds in any case.

Step 2.2: We now consider $\sigma = (a_i)_{i \in I}$ in \mathfrak{S}_0 with $|\sigma| = t + 1$ and representation $\psi(\sigma) = ((p_1, m_1), \dots, (p_q, m_q))$, $q \geq 1$. Let $T_1 = \omega \rho_1^\alpha \rho_2^\beta \rho_3$ and consider the rule $R' = T_2 T_3 \dots$ (which may have only one term, in which case Step 1 of the proof must be

used). Observe that since $p_1 > 0$ and $m_1 > 0$, ω does not modify p_1, m_1 and ρ_1, ρ_2 leave $p_1 > 0, m_1 > 0$, the action of ρ_3 will decrease p_1 and m_1 by at least one unit each. It then follows that $|T_1(\sigma)| \leq |\sigma| - 2 = t - 1$.

Moreover, by definition we have $R(\sigma) = R'(T_1(\sigma))$ and so $J_\sigma^R = J_\sigma^{T_1} \cup J_{T_1(\sigma)}^{R'}$. Applying induction hypothesis, we find

$$\varepsilon = R'((T_1(\sigma))_{J_{T_1(\sigma)}^{R'}}).$$

We claim that $R'((T_1(\sigma))_{J_{T_1(\sigma)}^{R'}}) = R(\sigma_{J_\sigma^R})$, which finishes the proof.

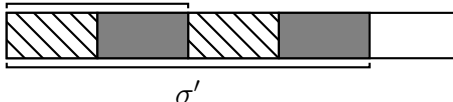
PROOF OF THE CLAIM: We have by definition

$$R(\sigma_{J_\sigma^R}) = R'(T_1(\sigma_{J_\sigma^R})) = R'(T_1(\sigma_{J_\sigma^{T_1}}, \sigma_{J_{T_1(\sigma)}^{R'}}))$$


It remains then to prove

$$T_1(\sigma_{J_\sigma^{T_1}}, \sigma_{J_{T_1(\sigma)}^{R'}}) = (T_1(\sigma))_{J_{T_1(\sigma)}^{R'}}. \quad (7)$$


We distinguish several cases on the structure of $R(\sigma)$, distinguishing the effect of T_1 and R' . In the following figures, the hatched and the grey parts are respectively the effect of T_1 and R' . The second hatched block may be empty (if $\omega = \varepsilon$), and similarly for the second grey block. We have to show that applying T_1 to the sequence formed by the hatched and grey parts (called σ') deletes exactly the hatched part. The foregoing analysis shows that only three cases can happen.

- Case 1: 

Observe that the action of ω is on the second hatched part, and depends only on (p_1, m_1) and the current (p_k, m_k) and not on subsequent values. As these pairs are unchanged, the action of ω on σ' or on σ is exactly the same. Now, ρ acts only on (p_1, m_1) and depends only on (p_1, m_1) . Hence, the result on σ and σ' is again the same.

- Case 2: 

Here (p_1, m_1) is changed into (p'_1, m'_1) in σ' , but both p'_1, m'_1 are positive due to the presence of ρ_3 in T_1 . Hence the same reasoning as in case 1 applies.

- Case 3: 

This case is treated like case 1.

(v) Clear from the foregoing analysis with head and tail.

□

3 Weak \mathbb{Z} -means based on the symmetric maximum

Specifically, we say that $\mu : \tilde{C}^* \rightarrow \tilde{C}$ is a *weak \mathbb{Z} -mean* if it satisfies anonymity, internality, monotonicity and *restricted decomposability*: for every sequence $(a_i)_{i \in I}$, there exists at least one nonempty $K \subseteq I$ on which decomposability applies:

$$\mu(a_K) = b \text{ implies } \mu(a_K, a_{I \setminus K}) = \mu(|K| \cdot b, a_{I \setminus K}),$$

where $k \cdot b$, for some nonnegative integer k and some $b \in \tilde{C}$, means b, b, \dots, b (repeated k times). In this case we say that μ is *decomposable for K* .

We recall that an anonymous function $\mu : \tilde{C}^* \rightarrow \tilde{C}$ is *associative* if $\mu(a) = a$ for every $a \in \tilde{C}$, and for every sequence $\sigma = (a_i)_{i \in I}$ and $K \subset I$, we have $\mu(\sigma) = \mu(\mu(a_K), a_{I \setminus K})$ (see [4, §2.3.1]). We say that an anonymous function μ is *weakly associative* if associativity holds for associative sequences (i.e., on $\mathfrak{S} \setminus \mathfrak{S}_0$).

We show in this section that the symmetric maximum is the basis for an infinite (uncountable) family of weakly associative weak \mathbb{Z} -means, namely $(\mathbb{O}_R)_{R \in \mathfrak{R}}$.

Lemma 2. Let $R \in \mathfrak{R}$.

- (i) For every sequence $\sigma = (a_i)_{i \in I}$, \mathbb{O}_R satisfies decomposability for every $K \supseteq J_\sigma^R$ and for every nonempty $K' \subseteq I \setminus J_\sigma^R$;
- (ii) For each rule R , there exists a sequence $\sigma = (a_i)_{i \in I}$ in \mathfrak{S}_0 and a set $K \subseteq I$ such that \mathbb{O}_R is not decomposable for K .

Proof. (i) We know by Lemma 1(iv) that $\mathbb{O}_R(a_{J_\sigma^R}) = 0$. Then, by Lemma 1(i) and the fact that 0 is neutral,

$$\mathbb{O}_R(|J_\sigma^R| \cdot 0, a_{I \setminus J_\sigma^R}) = \mathbb{O}_R(a_{I \setminus J_\sigma^R}) = \mathbb{O}(\sigma),$$

hence decomposability holds for J_σ^R .

Let $K \supset J_\sigma^R$ and let $b' = \mathbb{O}_R(a_K)$. We have to prove that

$$\mathbb{O}_R(|K| \cdot b', a_{I \setminus K}) = \mathbb{O}_R(\sigma) =: b.$$

Suppose that $R(\sigma)$ has a head, and hence b is in the head. Then $J_{a_K}^R = J_\sigma^R$ and $|b'| \leq |b|$. Moreover $b' = b$ holds iff K intersects the head of σ , otherwise strict inequality occurs and b' is the extremal value of the tail. Then $(|K| \cdot b', a_{I \setminus K})$ is associative in any case, because the head does not contain $-b$ and the tail is associative. Its extremal value is b in any case, because either K contains the head, and then $b' = b$, otherwise it does not and b is present in $a_{I \setminus K}$.

Suppose now that $R(\sigma)$ has no head. Observe that the sequence restricted to $K \cap I \setminus J_\sigma^R$ is not necessarily associative, therefore $J_{a_K}^R \supseteq J_\sigma^R$, but anyway $|b'| \leq |b|$. If $|b'| = |b|$, then the sequence restricted to $K \cap I \setminus J_\sigma^R$ is associative and $b' = b$ because no $-b$ is present in the tail. Then $J_{a_K}^R = J_\sigma^R$ and $\mathbb{O}_R(|K| \cdot b, a_{I \setminus K}) = b$. Suppose on the contrary that $|b'| < |b|$. Then the sequence $(|K| \cdot b', a_{I \setminus K})$ is associative because b is extremal and no $-b$ is present in the tail, therefore $\mathbb{O}_R(|K| \cdot b', a_{I \setminus K}) = b$.

Consider finally $K' \subseteq I \setminus J_\sigma^R$ and let $b' = \mathbb{O}_R(a_{K'})$.

Suppose that $R(\sigma)$ has a head, and that K intersects the head. Then clearly $b' = b$ and $\mathbb{O}_R(|K'| \cdot b, a_{I \setminus K'}) = b$, as desired. If on the contrary K' is included in the tail, then $|b'| < |b|$. By Lemma 1(iii), $\mathbb{O}_R(|K'| \cdot b', a_{I \setminus K'}) = b$.

Suppose now that $R(\sigma)$ has no head, which implies that K' is included in the tail. Then $|b'| \leq |b|$. If equality occurs, then $b' = b$, $a_{K'}$ is associative and the modified tail remains associative. It follows that the desired result holds by Lemma 1(iii). If $|b'| < |b|$, the tail remains also associative because there is no $-b$. Again, the result holds by Lemma 1(iii).

(ii) We prove now that for each rule, one can find a nonassociative sequence and a subsequence of it on which the symmetric maximum is not decomposable. Consider the sequence $\sigma = (2, -2, 1)$. Suppose $R \in \mathfrak{R}_{123}$ and take the subsequence $\sigma_K = (-2, 1)$. Then $\mathbb{O}_R(-2, 1) = -2$ and $\mathbb{O}_R(\sigma) = 1$. However, $\mathbb{O}_R(2, -2, -2)$ is equal to either 0 or -2 , depending whether R uses ρ_2 or not. Hence decomposability fails for K .

Suppose now that $R \notin \mathfrak{R}_{123}$ and consider the subsequence $\sigma_K = (2, -2)$. Then $\mathbb{O}_R(2, -2) = 0$ and $\mathbb{O}_R(\sigma) = 0$ (supposing that R contains ρ_4 . Otherwise, replace σ by $(2, -2, -1)$). However, $\mathbb{O}_R(0, 0, 1) = 1$, and decomposability again fails on K . \square

Corollary 1. (i) The symmetric maximum is decomposable for no computation rule.

(ii) The symmetric maximum is decomposable for any sequence in $\mathfrak{S} \setminus \mathfrak{S}_0$, for any computation rule.

Observe also that $\langle \cdot \rangle_-^+$ is the only rule yielding an extremal operator. Indeed,

$$\mathbb{O}(\langle (a_i)_{i \in I} \rangle_-^+) = \left(\min_{i \in I} a_i \right) \mathbb{O} \left(\max_{i \in I} a_i \right).$$

Any other rule (as shown in [2]) deletes less terms than $\langle \cdot \rangle_-^+$, hence there are sequences for which other elements than the minimum and maximum are used. This fact motivates the definition of a weaker notion of extremality (see below).

For characterization purposes, we introduce a number of properties.

Definition 1. Let $\mu : \tilde{C}^* \rightarrow \tilde{C}$.

(i) μ is *extremal w.r.t.* $R \in \mathfrak{R}$ if for every $\sigma = (a_i)_{i \in I}$ in \mathfrak{S}

$$\mu(\sigma) = \left(\min_{i \in I \setminus J_\sigma^R} a_i \right) \mathbb{O} \left(\max_{i \in I \setminus J_\sigma^R} a_i \right);$$

(ii) μ satisfies the *replacement property w.r.t.* $R \in \mathfrak{R}$, abbreviated by $\text{RP}(R)$, if for every sequence $\sigma = (a_i)_{i \in I}$ in \mathfrak{S} ,

$$\mu(\sigma) = b \Rightarrow \begin{cases} \mu(|J_\sigma^R| \cdot b, a_{I \setminus J_\sigma^R}) = b \\ \mu(a_{J_\sigma^R}, |I \setminus J_\sigma^R| \cdot b) = b \end{cases}, \text{ and}$$

(iii) μ is *retractive w.r.t.* $R \in \mathfrak{R}$, if for every $\sigma = (a_i)_{i \in I} \in \mathfrak{S}$ we have that $\mu(a_{I \setminus J_\sigma^R}) = \mu(\sigma)$;

- (iv) μ is *weakly retractive w.r.t.* $R \in \mathfrak{R}$ if for every $\sigma = (a_i)_{i \in I} \in \mathfrak{S}$, we have that $|\mu(\sigma)| \geq |\mu(a_{I \setminus J_\sigma^R})|$ with both of the same sign, i.e.,

$$\mu(\sigma) \geq \mu(a_{I \setminus J_\sigma^R}) \geq 0 \quad \text{or} \quad \mu(\sigma) \leq \mu(a_{I \setminus J_\sigma^R}) \leq 0.$$

Lemma 3. For any $R \in \mathfrak{R}$, the symmetric maximum \mathbb{O}_R is extremal, (weakly) retractive w.r.t. R , and satisfies $\text{RP}(R)$.

Proof. Retractiveness and extremality are established in Lemma 1(i). As for $\text{RP}(R)$, $b = \mathbb{O}_R(\sigma) = \mathbb{O}_R(a_{I \setminus J_\sigma^R})$ by retractiveness, and $\mathbb{O}_R(a_{J_\sigma^R}, |I \setminus J_\sigma^R| \cdot b) = b$ by decomposability for $K = I \setminus J_\sigma^R$. Now, since $a_{I \setminus J_\sigma^R}$ is associative and its extremal term is b , the sequence $(|J_\sigma^R| \cdot b, a_{I \setminus J_\sigma^R})$ is still associative and b is its extremal term. \square

Remark 1. For any $R \in \mathfrak{R}$, the $\text{RP}(R)$ property implies *unanimity*, i.e., $\mu(b, \dots, b) = b$, for every $b \in \text{Ran}(\mu(\mathfrak{S} \setminus \mathfrak{S}_0))$, where $\text{Ran}(\cdot)$ indicates the range of a function. Indeed, for $\sigma \in \mathfrak{S} \setminus \mathfrak{S}_0$ with $\mu(\sigma) = b$, $J_\sigma^R = \emptyset$.

4 Characterizations

We give several characterizations of weak \mathbb{Z} -means based on the symmetric maximum. The first one involves decomposability. We define $\text{Ker}(\mu) = \{\sigma \in \mathfrak{S}^* \mid \mu(\sigma) = 0\}$, where \mathfrak{S}^* is the set of sequences in \mathfrak{S} not containing 0. Observe that by our convention, $\text{Ker}(R) = \text{Ker}(\mathbb{O}_R)$.

4.1 Characterization by decomposability

Theorem 3. Let $\mu : \tilde{C}^* \rightarrow \tilde{C}$ and let $R \in \mathfrak{R}$. Then $\mu = \mathbb{O}_R$ if and only if μ satisfies (C1), (C2), (C3) on \tilde{C}^2 , is anonymous, weakly associative, decomposable on J_σ^R for every sequence σ , and satisfies $\mu(\sigma) = 0$ for every $\sigma \in \text{Ker}(R)$ (i.e., $\text{Ker}(R) \subseteq \text{Ker}(\mu)$).

The following observation is useful in many subsequent proofs.

Observation 1. Let $\mu : \tilde{C}^* \rightarrow \tilde{C}$. If μ satisfies (C1), (C2), (C3), and weak associativity, then $\mu = \mathbb{O}$ on $\mathfrak{S} \setminus \mathfrak{S}_0$.

Indeed, by Theorem 2 (C1), (C2), (C3) and weak associativity imply that μ coincides with \mathbb{O} on \tilde{C}^2 . It thus follows by weak associativity that $\mu(\sigma) = \mathbb{O}(\sigma)$ for every associative sequence σ .

Proof. The “only if” part is obtained from previous results (in particular Lemma 2).

For the “if” part, by Observation 1, $\mu = \mathbb{O}$ on $\mathfrak{S} \setminus \mathfrak{S}_0$.

Consider now sequences in \mathfrak{S}_0 and, for the sake of a contradiction, suppose that $\mu \neq \mathbb{O}_R$ on \mathfrak{S}_0 . Take $\sigma = (a_i)_{i \in I}$ in \mathfrak{S}_0 such that $\mu(\sigma) \neq \mathbb{O}_R(\sigma)$ and $|\text{Ran}(\sigma)|$ is minimum, where $\text{Ran}(\sigma)$ is the range of σ (i.e., the set of all integers in σ).

Suppose first that $|\text{Ran}(\sigma)| \geq 3$. Note that $|\text{Ran}(a_{J_\sigma^R})| \geq 2$, so that $|\text{Ran}(|J_\sigma^R| \cdot \mu(a_{J_\sigma^R}), a_{I \setminus J_\sigma^R})| < |\text{Ran}(\sigma)|$. By definition of σ and by applying decomposability on J_σ^R , it follows that

$$\mu(\sigma) = \mu(a_{J_\sigma^R}, a_{I \setminus J_\sigma^R}) = \mu(|J_\sigma^R| \cdot \mu(a_{J_\sigma^R}), a_{I \setminus J_\sigma^R}) = \mathbb{O}_R(|J_\sigma^R| \cdot \mu(a_{J_\sigma^R}), a_{I \setminus J_\sigma^R}). \quad (8)$$

Now, by Lemma 1(iv), $a_{J_\sigma^R} \in \text{Ker}(R) \subseteq \text{Ker}(\mu)$, and since 0 is neutral for \mathbb{O}_R , we get

$$\begin{aligned}\mathbb{O}_R(|J_\sigma^R| \cdot \mu(a_{J_\sigma^R}), a_{I \setminus J_\sigma^R}) &= \mathbb{O}_R(|J_\sigma^R| \cdot 0, a_{I \setminus J_\sigma^R}) \\ &= \mathbb{O}_R(a_{I \setminus J_\sigma^R}) = \mathbb{O}_R(\sigma)\end{aligned}$$

where the last equality comes from Lemma 1(i). Combining the latter result with (8) yields $\mu(\sigma) = \mathbb{O}_R(\sigma)$, a contradiction.

Suppose now that $|\text{Ran}(\sigma)| = 2$ (note that a lower value is impossible). Then necessarily σ is such that $\psi(\sigma) = (p_1, m_1)$ with $p_1 > 0$, $m_1 > 0$. Two cases arise. Suppose first that $a_{J_\sigma^R} = \sigma$. Then $\sigma \in \text{Ker}(R)$, hence $\mu(\sigma) = 0 = \mathbb{O}_R(\sigma)$, a contradiction. Suppose on the contrary that $a_{J_\sigma^R} \neq \sigma$. Since $J_\sigma^R \in \text{Ker}(R) \subseteq \text{Ker}(\mu)$, it follows from decomposability for J_σ^R that

$$\mu(\sigma) = \mu(a_{I \setminus J_\sigma^R}, |J_\sigma^R| \cdot \mu(a_{J_\sigma^R})) = \mu(a_{I \setminus J_\sigma^R}, |J_\sigma^R| \cdot 0). \quad (9)$$

Observe that by the assumption $|\text{Ran}(\sigma)| = 2$, we have necessarily $|\text{Ran}(a_{I \setminus J_\sigma^R})| = 1$. Hence the sequence $(a_{I \setminus J_\sigma^R}, |J_\sigma^R| \cdot 0)$ is associative, so that

$$\mu(a_{I \setminus J_\sigma^R}, |J_\sigma^R| \cdot 0) = \mathbb{O}_R(a_{I \setminus J_\sigma^R}, |J_\sigma^R| \cdot 0) = \mathbb{O}_R(a_{I \setminus J_\sigma^R}) = \mathbb{O}_R(\sigma)$$

Combining this with (9), we find $\mu(\sigma) = \mathbb{O}_R(\sigma)$, again a contradiction. \square

Remark 2. If $R = \langle \cdot \rangle_-^+$, then decomposability on J_σ^R is trivially satisfied and therefore is not needed in the axiomatization. Indeed, $J_\sigma^{\langle \cdot \rangle_-^+} = I$ for any $\sigma = (a_i)_{i \in I} \in \mathfrak{S}_0$, and letting $\mu(\sigma) = b$, we have $\mu(b, \dots, b) = \mathbb{O}(b, \dots, b) = b$ by Observation 1.

4.2 Characterization by the replacement property and others

We introduce further characterizations based on the properties in Definition 1. First, we observe that the kernels of μ and R are somehow related.

Lemma 4. Suppose that $\mu : \tilde{C}^* \rightarrow \tilde{C}$ satisfies (C1), (C2), (C3) on \tilde{C}^2 and weak associativity. If μ verifies $\text{RP}(R)$ for some $R \in \mathfrak{R}$, then $\text{Ker}(\mu) \subseteq \text{Ker}(R)$. If in addition $|\mathbb{O}_R(\sigma)| \geq |\mu(\sigma)|$ for every $\sigma \in \mathfrak{S}$, denoted by $|\mathbb{O}_R| \geq |\mu|$, then $\text{Ker}(R) = \text{Ker}(\mu)$.

Proof. Note that by (C2) $\text{Ker}(\mu) \neq \emptyset$. For the sake of a contradiction, suppose that there exists $\sigma = (a_i)_{i \in I} \in \text{Ker}(\mu) \setminus \text{Ker}(R)$, i.e., $\mu(\sigma) = 0$ and $\mathbb{O}_R(\sigma) \neq 0$. As before, take such a σ with minimum $|\text{Ran}(\sigma)|$. As μ satisfies (C1) and (C2), we may assume that $|\text{Ran}(\sigma)| \geq 2$.

The case $|\text{Ran}(\sigma)| = 2$ is dealt with as above. So suppose that $|\text{Ran}(\sigma)| \geq 3$. Observe that $J_\sigma^R = \emptyset$ is impossible. Indeed, this would mean that $\sigma \in \mathfrak{S}^* \setminus \mathfrak{S}_0$, and by Observation 1, $0 = \mu(\sigma) = \mathbb{O}(\sigma)$, a contradiction. Hence, $J_\sigma^R \neq \emptyset$. Then

$$|\text{Ran}(|J_\sigma^R| \cdot 0, a_{I \setminus J_\sigma^R})| < |\text{Ran}(\sigma)|. \quad (10)$$

By $\text{RP}(R)$,

$$\mu(a_{J_\sigma^R}, |I \setminus J_\sigma^R| \cdot 0) = 0 = \mu(|J_\sigma^R| \cdot 0, a_{I \setminus J_\sigma^R}).$$

By the minimality assumption on σ , it then follows from (10) that

$$0 = \mu(|J_\sigma^R| \cdot 0, a_{I \setminus J_\sigma^R}) = \mathbb{O}_R(|J_\sigma^R| \cdot 0, a_{I \setminus J_\sigma^R}).$$

By Lemma 1 (i) and the fact that 0 is neutral in \mathbb{O}_R , we then have that

$$0 = \mathbb{O}_R(|J_\sigma^R| \cdot 0, a_{I \setminus J_\sigma^R}) = \mathbb{O}_R(a_{I \setminus J_\sigma^R}) = \mathbb{O}_R(\sigma),$$

which contradicts the choice of σ . Hence there is no $\sigma \in \text{Ker}(\mu) \setminus \text{Ker}(R)$, and thus $\text{Ker}(\mu) \subseteq \text{Ker}(R)$.

The last assertion is obvious. \square

Remark 3. If $\mu : \tilde{C}^* \rightarrow \tilde{C}$ satisfies (C1), (C2), (C3), is weakly associative and weakly retractive w.r.t. R , then by Observation 1, $|\mathbb{O}_R| \leq |\mu|$.

Proposition 1. Let $\mu : \tilde{C}^* \rightarrow \tilde{C}$ be a function that satisfies (C1), (C2), (C3) on \tilde{C}^2 . Then the following assertions are equivalent:

- (i) $\mu = \mathbb{O}_R$;
- (ii) μ is extremal w.r.t. $R \in \mathfrak{R}$;
- (iii) μ is weakly associative and it is retractive w.r.t. R ;
- (iv) μ is weakly associative, weakly retractive, and $|\mathbb{O}_R| \geq |\mu|$.

Proof. By Lemma 1(i), (i) is equivalent to (ii) and (i) implies (iii). Also, if μ satisfies (C1), (C2), (C3) on \tilde{C}^2 and it is weakly associative and retractive w.r.t. R , then

$$\mu(\sigma) = \mu(a_{I \setminus J_\sigma^R}) = \left(\min_{i \in I \setminus J_\sigma^R} a_i \right) \mathbb{O} \left(\max_{i \in I \setminus J_\sigma^R} a_i \right),$$

that is, μ is extremal w.r.t. R , and hence (iii) implies (ii).

Also, by Lemma 3 we have that (i) implies (iv). To show that the converse implication also holds, note that by Remark 3 and $|\mathbb{O}_R| \geq |\mu|$ we get $|\mathbb{O}_R| = |\mu|$ and, by weak retractiveness and Observation 1, it follows that $\mathbb{O}_R = \mu$. \square

Lemma 5. Suppose that $\mu : \tilde{C}^* \rightarrow \tilde{C}$ satisfies (C1), (C2), (C3) on \tilde{C}^2 , and that μ is weakly associative, anonymous, and satisfies $\text{RP}(R)$ and $|\mathbb{O}_R| \geq |\mu|$ with both of the same sign (i.e., $\mathbb{O}_R(\sigma) \geq 0$ if and only if $\mu(\sigma) \geq 0$), for some $R \in \mathfrak{R}$. Then $\mu = \mathbb{O}_R$.

Proof. First we show that $|\mathbb{O}_R| = |\mu|$. To this end, note that from (C1), (C2), (C3) and weak associativity it follows by Observation 1 that for any sequence $\sigma = (a_i)_{i \in I} \in \mathfrak{S}_0$,

$$\mu(a_{I \setminus J_\sigma^R}) = \mathbb{O}_R(a_{I \setminus J_\sigma^R}) = \mathbb{O}_R(\sigma) \quad (11)$$

and, since $|\mathbb{O}_R| \geq |\mu|$, we cannot have $|\mu(\sigma)| > |\mu(a_{I \setminus J_\sigma^R})|$. Hence, $|\mu(\sigma)| \leq |\mu(a_{I \setminus J_\sigma^R})|$.

Now, suppose that $|\mu(\sigma)| < |\mu(a_{I \setminus J_\sigma^R})|$. By (11) and Lemma 1(i),

$$|\mu(a_{I \setminus J_\sigma^R})| = |\mathbb{O}(a_{I \setminus J_\sigma^R})| = |\mathbb{O}_R(a_{I \setminus J_\sigma^R})| = |\mathbb{O}_R(\sigma)| = \max_{i \in I \setminus J_\sigma^R} |a_i|. \quad (12)$$

Let $b = \mu(\sigma)$. Since μ and \mathbb{O}_R have the same sign and μ is odd, we may assume that $b \geq 0$. By assumption and (12), we have that

$$b < \max_{i \in I \setminus J_\sigma^R} |a_i| = |\mathbb{O}_R(\sigma)|. \quad (13)$$

By $RP(R)$, we have that $\mu(\sigma) = \mu(|J_\sigma^R| \cdot b, a_{I \setminus J_\sigma^R})$, and since $(|J_\sigma^R| \cdot b, a_{I \setminus J_\sigma^R})$ is associative (by (13)),

$$\mu(\sigma) = \mu(|J_\sigma^R| \cdot b, a_{I \setminus J_\sigma^R}) = \mathbb{O}_R(|J_\sigma^R| \cdot b, a_{I \setminus J_\sigma^R}). \quad (14)$$

However, \mathbb{O}_R is extremal w.r.t. R and from the fact that $0 \leq b < \max_{i \in I \setminus J_\sigma^R} |a_i|$ it follows from Lemma 1(i) that

$$\mathbb{O}_R(|J_\sigma^R| \cdot b, a_{I \setminus J_\sigma^R}) = \mathbb{O}_R(a_{I \setminus J_\sigma^R}) = \mathbb{O}_R(\sigma),$$

and thus by (14)

$$\mu(\sigma) = \mu(|J_\sigma^R| \cdot b, a_{I \setminus J_\sigma^R}) = \mathbb{O}_R(|J_\sigma^R| \cdot b, a_{I \setminus J_\sigma^R}) = \mathbb{O}_R(\sigma).$$

Hence, $\mathbb{O}_R(\sigma) = b$, which contradicts (13), and thus $|\mu(\sigma)| = |\mu(a_{I \setminus J_\sigma^R})| = |\mathbb{O}_R(\sigma)|$, as desired. Since \mathbb{O}_R and μ have the same sign, it then follows that $\mathbb{O}_R = \mu$. \square

Remark 4. As for the previous characterization, the rule $R = \langle \cdot \rangle_-^+$ is particular, because the property $|\mathbb{O}_R| \geq |\mu|$ with the same sign is enough to characterize it, without the help of $RP(R)$.

We can now state further characterizations as mentioned earlier.

Theorem 4. Let $\mu : \tilde{C}^* \rightarrow \tilde{C}$ be an anonymous and weakly associative function that satisfies (C1), (C2), (C3) on \tilde{C}^2 . For every $R \in \mathfrak{R}$, the following assertions are equivalent.

- (i) $\mu = \mathbb{O}_R$.
- (ii) μ decomposable on J_σ^R for every sequence σ , and $\text{Ker}(R) \subseteq \text{Ker}(\mu)$.
- (iii) μ is extremal w.r.t. R .
- (iv) μ is retractive w.r.t. R .
- (v) μ is weakly retractive and $|\mathbb{O}_R| \geq |\mu|$.
- (vi) μ verifies $RP(R)$ and $|\mathbb{O}_R| \geq |\mu|$ both with the same sign.

Proof. By Theorem 3 and Proposition 1 we have that (i) to (v) are all equivalent. By Lemma 5 we also have that (vi) implies (i), and by Lemma 3 it follows that (i) implies (vi). \square

4.3 Independence of the axioms

As a preliminary remark, we must distinguish the roles of the different axioms introduced. Axioms (C1), (C2), (C3) pertain only to the case of a binary operator and are clearly independent. Weak associativity permits to extend to associative sequences, and by Observation 1, the operator \mathbb{O} is obtained. Hence, the four axioms are independent.

The remaining axioms only serve to determine what happens for nonassociative sequences. Anonymity permits to work on ordered sequences and is clearly independent of the other remaining axioms.

Therefore, our task consists in proving that, for any rule R (except for $R = \langle \cdot \rangle_-^+$, see Remarks 2 and 4)

- (i) in Theorem 3, decomposability on J_σ^R and $\text{Ker}(R) \subseteq \text{Ker}(\mu)$ are independent, and
- (ii) in Lemma 5, $\text{RP}(R)$ and $|\mathbb{Q}_R| \geq |\mu|$ with the same sign, are also independent.

We begin with Theorem 3 and suppose that μ satisfies (C1), (C2), (C3), anonymity, and weak associativity. Consider any rule $R \neq \langle \cdot \rangle_-^+$ in \mathfrak{R} . Clearly, we only have to consider nonassociative sequences $\sigma = (a_i)_{i \in I}$. Let $C = \mathbb{Z}$.

- (i) Define $\mu(\sigma)$ as $\min(a_{I \setminus J_\sigma^R})$ for any nonassociative σ s.t. $I \setminus J_\sigma^R \neq \emptyset$, and 0 otherwise. Then μ satisfies $\text{Ker}(R) \subseteq \text{Ker}(\mu)$ but not decomposability. Indeed, by Lemma 1(iv)

$$\mu(a_{J_\sigma^R}) = 0 \text{ and } \mu(|J_\sigma^R| \cdot 0, a_{I \setminus J_\sigma^R}) = \mathbb{Q}(|J_\sigma^R| \cdot 0, a_{I \setminus J_\sigma^R}) = \mathbb{Q}(a_{I \setminus J_\sigma^R}) \neq \min(a_{I \setminus J_\sigma^R})$$

in general. For instance, take $\sigma = (3, -3, 2, 1)$.

- (ii) Define $\mu(a_{J_\sigma^R}) = \alpha \in C$, $\alpha > 0$, for any σ , except if $a_{J_\sigma^R} = (a, -a)$, in which case $\mu(a_{J_\sigma^R}) = 0$, so that (C2) is satisfied. Then, $\text{Ker}(R) \subseteq \text{Ker}(\mu)$ is not satisfied. We claim that, assuming that μ is decomposable on J_σ^R , μ is well defined on \mathfrak{S}_0 (note that by (C2), we already have $\mu(a, -a) = 0$ for every a). Indeed, by decomposability, $\mu(\sigma) = \mu(\sigma_{J_\sigma^R}, \sigma_{I \setminus J_\sigma^R}) = \mu(|J_\sigma^R| \cdot \alpha, \sigma_{I \setminus J_\sigma^R})$ or $\mu(|J_\sigma^R| \cdot 0, \sigma_{I \setminus J_\sigma^R})$. In the latter case, $(|J_\sigma^R| \cdot 0, \sigma_{I \setminus J_\sigma^R}) =: \sigma'$ is associative, so that $\mu(\sigma) = \mathbb{Q}(\sigma')$. In the former case, if $-\alpha$ is not extremal in $\sigma_{I \setminus J_\sigma^R}$, the sequence $\sigma' = (|J_\sigma^R| \cdot \alpha, \sigma_{I \setminus J_\sigma^R})$ is associative, and therefore $\mu(\sigma) = \mathbb{Q}(\sigma')$. Otherwise, $\sigma_{I \setminus J_\sigma^R}$ contains $-\alpha$ as extremal term and possibly other terms smaller in magnitude. By repeatedly applying decomposability, say $k \geq 1$ times, we can delete the occurrence of $-\alpha$ in the subsequence that remains undeleted, and thus we arrive at a sequence $\sigma^{(k)} = (|J_{\sigma^{(k-1)}}^R| \cdot \alpha, \sigma_{I \setminus J_{\sigma^{(k-1)}}^R}^{(k-1)})$ that is associative. To illustrate this fact, suppose that $\sigma = (2 \cdot 3, 7 \cdot (-3))$, $R = \langle \cdot \rangle_0 = \rho_3^*$, and that $\mu(n \cdot 3, n \cdot (-3)) = \alpha = 3$, for $n \geq 2$. Then the first application of decomposability produces $\sigma' = (|J_\sigma^R| \cdot \alpha, \sigma_{I \setminus J_\sigma^R}) = (4 \cdot 3, 5 \cdot (-3))$ and thus $\sigma'_{I \setminus J_{\sigma'}^R} = (-3)$. By reapplying decomposability, we obtain $\sigma^{(2)} = (|J_{\sigma'}^R| \cdot \alpha, \sigma'_{I \setminus J_{\sigma'}^R}) = (8 \cdot 3, 1 \cdot (-3))$ and in this case we have $\sigma_{I \setminus J_{\sigma^{(2)}}^R}^{(2)} = (7 \cdot 3)$. By applying decomposability for a third time, we get $\sigma^{(3)} = (|J_{\sigma^{(2)}}^R| \cdot \alpha, \sigma_{I \setminus J_{\sigma^{(2)}}^R}^{(2)}) = (9 \cdot 3)$, which is now associative, and from which it follows that

$$\mu(\sigma) = \mu(\sigma') = \mu(\sigma^{(2)}) = \mu(\sigma^{(3)}) = \mathbb{Q}(\sigma^{(3)}) = 3.$$

In general, if $k \geq 1$ applications of decomposability suffice to delete the occurrences of $-\alpha$ in $\sigma_{I \setminus J_\sigma^R}$ (i.e., $\sigma^{(k)}$ is associative), then

$$\mu(\sigma) = \mu(\sigma') = \dots = \mu(\sigma^{(k)}) = \mathbb{Q}(\sigma^{(k)}) = \alpha.$$

Let us now turn to Lemma 5 and suppose as before that $R \neq \langle \cdot \rangle_-^+$.

- (i) Define μ by $\mu(\sigma) = 0$ for every $\sigma \in \text{Ker}(R)$, otherwise, $\mu(\sigma)$ is the value in $\sigma_{I \setminus J_\sigma^R}$ which has smallest absolute value, with the sign of the extremal value of $\sigma_{I \setminus J_\sigma^R}$ (Example: $\mu(3, -3, -2, 1) = -1$ with $R = \langle \cdot \rangle_0$). Then μ satisfies $|\mathbb{Q}_R| \geq |\mu|$ with the same sign. However, $\text{RP}(R)$ is not satisfied, whenever σ is such that $I \setminus J_\sigma^R \neq \emptyset$, with at least two values of different magnitude.

(ii) Consider $\tilde{C} = \mathbb{Z}$ and define μ by $\mu(\sigma) = 0$ for every $\sigma \in \text{Ker}(R)$, otherwise

$$\mu(\sigma) = \text{sgn}(\text{extremal value of } \sigma_{I \setminus J_\sigma^R}) \cdot (\min_{i \in J_\sigma^R}(|a_i|) - 1).$$

For example, $\mu(4, -4, 3, -3, -1) = -2$ with $R = \langle \cdot \rangle_0$. As the example shows, the property $|\otimes_R| \geq |\mu|$ fails. We claim that μ satisfies $RP(R)$. Indeed, for any $\sigma \notin \text{Ker}(R)$, letting $\mu(\sigma) = b$, we have $\mu(|J_\sigma^R| \cdot b, \sigma_{I \setminus J_\sigma^R}) = b$ because $(|J_\sigma^R| \cdot b, \sigma_{I \setminus J_\sigma^R})$ is associative with extremal value b , and $\mu(\sigma_{J_\sigma^R}, |I \setminus J_\sigma^R| \cdot b) = b$, since for $\sigma' = (\sigma_{J_\sigma^R}, |I \setminus J_\sigma^R| \cdot b)$ we have $J_{\sigma'}^R = J_\sigma^R$.

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